

High-SNR Analysis of Outage-Limited Communications With Bursty and Delay-Limited Information

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Abstract—This work analyzes the high-SNR asymptotic error performance of outage-limited communications with fading, where the number of bits that arrive at the transmitter during any timeslot is random but the delivery of bits at the receiver must adhere to a strict delay limitation. Specifically, bit errors are caused by erroneous decoding at the receiver or violation of the strict delay constraint. Under certain scaling of the statistics of the bit-arrival process with SNR, this paper shows that the optimal decay behavior of the asymptotic total probability of bit error depends on how fast the burstiness of the source scales down with SNR. If the source burstiness scales down too slowly, the total probability of error is asymptotically dominated by delay-violation events. On the other hand, if the source burstiness scales down too quickly, the total probability of error is asymptotically dominated by channel-error events. However, at the proper scaling, where the burstiness scales linearly with $\frac{1}{\sqrt{\log \text{SNR}}}$ and at the optimal coding duration and transmission rate, the occurrences of channel errors and delay-violation errors are asymptotically balanced. In this latter case, the optimal exponent of the total probability of error reveals a tradeoff that addresses the question of how much of the allowable time and rate should be used for gaining reliability over the channel and how much for accommodating the burstiness with delay constraints.

Index Terms—Batch service, delay effects, diversity, error analysis, fading channels, large deviations, queueing analysis.

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I. INTRODUCTION

THIS work analyzes the high signal-to-noise-ratio (SNR) performance of outage-limited communications where the information to be communicated is delay-limited and where the information arrives at the transmitter in a stochastic manner. We consider the following setting (Fig. 1) in our study.

- A random number of bits arrive at the transmitter during any given timeslot. Bits accumulate in an infinite buffer while waiting for their turn to be bunched into codewords and transmitted under a first-come, first-transmit policy.
- There is no feedback to the transmitter; retransmission of the bits in error is not considered.
- Communication over the fading channel is outage-limited ([1], [2]), where the transmitter is unaware of the instantaneous channel state and, as a consequence, operates at a fixed transmission rate, R . During a deep fade (also known as an outage), the channel seen by the decoder is too weak to allow recovery of the data content from the transmitted signal. Characteristic settings are those of multiple-input–multiple-output (MIMO) and cooperative outage-limited communications.
- Coding takes place in blocks where each codeword spans over a fixed and finite integral number, T , of timeslots. Each codeword has an information content of RT bits. In addition, coding is “fully-diverse,” i.e., the decoding at the receiver takes place only at the end of the coding block.
- The delay bound, D , is a maximum allowable time duration from the moment a bit arrives at the transmitter until the moment it is decoded at the receiver. The delay experienced by a bit is the sum of the time spent waiting in the buffer and the time spent in the block decoding process. Note that the waiting time in the buffer is random due to the stochastic arrival process.
- A bit is declared in error either when it is decoded incorrectly at the decoder, or when it violates the delay bound.

For the above setting, we are interested in the high-SNR asymptotic total probability of bit error. Note that for a given transmission rate, R , and a coding block duration, T , there exists a tradeoff between the probabilities of decoding error versus the delay violation. We expect that longer coding blocks allow the encoded bits to be transmitted over more fading realizations and hence, achieve higher diversity and fewer decoding errors. However, longer coding blocks cause more bits to violate the delay requirement. In other words, one intuitively expects that there is an optimal choice of the fixed transmission

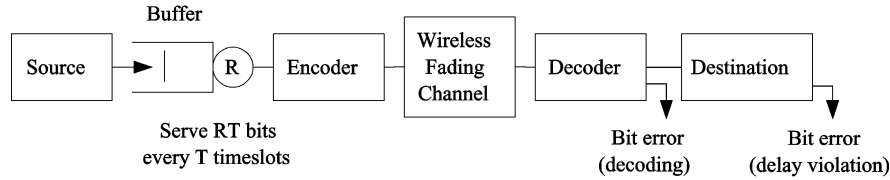


Fig. 1. System model.

rate, R , and the fixed coding block duration, T , for which the total probability of bit error is minimized. The goal of this paper is to identify these optimal quantities.

A. Prior Work and Our Contribution

High demands on the quality of service (QoS), in terms of both packet losses and packet delays, have fueled substantial research interest in jointly considering channels and queues. Communication of delay-sensitive bits over wireless channels has been addressed under various assumptions and settings in works such as [3]–[8]. Often, asymptotic approximations are employed to enable tractable analysis of the problem. Below we detail the existing work with their corresponding settings and the relationships to this paper.

The first group we discuss, [3]–[6], consists of scenarios where the current channel state information (CSI) is assumed to be known at both the transmitter and receiver. For example, in [3] and [4], Berry and Gallager address the tradeoff between the minimum average power consumption and the average delay (the power-delay tradeoff) over a Markovian fading channel with CSI both at the transmitter and the receiver. In such a setting, the transmitter dynamically varies power (i.e., the rate) in response to the current queue length and channel state. In [5], Rajan *et al.* derive optimal delay-bounded schedulers for transmission of constant-rate traffic over finite-state fading channels. In [6], Negi and Goel apply the effective capacity [9] and error exponent techniques to find the code-rate allocation that maximizes the decay rate of the asymptotic probability of error for a given asymptotically large delay requirement. Similar to [3] and [4], the proposed dynamic code-rate allocation in [6] is in response to the current channel fading and is possible by assuming CSI knowledge at the transmitter.

A second group of work (e.g., [7], [8]) focuses on scenarios where CSI is unknown to the transmitter but there is a mechanism for retransmission of codewords when the channel is in outage. As a tradeoff to protection against channel outage, this retransmission incurs extra delays to the bits in the buffer. In [7], for example, Bettesh and Shamaï (Shitz) address the problem of minimizing the average delay, under average power constraints and fixed transmission rate. They provide asymptotic analysis, under heavy load condition and asymptotically large queue length, for the optimal adaptive policies that adjust the transmission rate and/or transmission power in response to the current queue length at the transmitter. In another example, Liu *et al.* in [8] study the problem of optimal (fixed) transmission rate to maximize the decay rate of the probability of buffer

overflow for ON–OFF channels and Markov-modulated arrivals. The channel is considered “off” when outage occurs.

Although our work uses a similar performance measure to [6], namely the decay rate of the asymptotic probability of error, it covers the scenarios in which CSI is not available to the transmitter (no CSIT) and there is no retransmission. In such a setting, the variation of the fading channel is combatted via a coding over multiple independent fading realizations.¹ While this approach improves the transmission reliability, its longer coding duration increases the end-to-end delay any bit faces, and can potentially increase the probability of delay violation. In other words, in the absence of CSIT and retransmission, the transmission reliability, as well as the delay violation probability, are functions of the coding rate and duration. Consequently, our work compliments this previous research as it considers the effect of a delay violation requirement, in the absence of CSI at the transmitter and retransmission, on the operation of the physical layer. We consider a fixed transmission rate and code duration, as opposed to dynamic policies.

Since it is difficult to derive the exact relationship between the system parameters and the probabilities of channel decoding error and the delay violation, we choose to study an asymptotic approximation when the signal-to-noise ratio (SNR) is asymptotically high. The first advantage of this choice is the availability of an asymptotic high-SNR analysis for the channel decoding error probability. This high-SNR analysis is known as the *diversity-multiplexing-tradeoff* (DMT) analysis [1] and has received a great deal of attention during the past few years. Another advantage of the high-SNR analysis is that, for the class of arrival processes we consider in this paper, we can derive an asymptotic approximation of the delay violation probability that is valid even when the delay requirement D is finite and small. This derivation (Lemma 2) is based on a large-deviations result known as the Gärtner–Ellis theorem (see, e.g., [10]) and extends the large deviations exponent for a queue with asymptotic number of flows (as provided in [11]–[14]) to a queue with batch service discipline. Given that the asymptotic expression of the total probability of bit error is valid without requiring asymptotically large D , it is then meaningful to ask about the optimal coding block duration, a question which is not answered in studies with asymptotic D (e.g., [3], [4], [6], [15]–[17]).

We also would like to point out that our work was motivated by the work of Holliday and Goldsmith [18] where, under a high-SNR asymptotic approximation, the optimal operating channel transmission rate for a concatenated source/channel

¹For example, the multiple independent fading realizations can be a result of fading in multiple channel coherence time intervals (known as time diversity), or fading in multiple independent spatial channels, as in MIMO channel (spatial diversity), or cooperative relay channel (cooperative diversity).

system is studied. Following the approach in [18], we study a concatenated queue/channel system under a high-SNR approximation.

B. Overview of the Results

This work focuses on the notion of SNR error exponent as a measure of performance. Specifically, we are interested in finding how the asymptotic total probability of error decays with SNR. To keep the problem meaningful, we consider a scenario under which the overall traffic loading of the system (the ratio between the mean arrival rate and the ergodic capacity of the channel) is kept independent of SNR. That is, we consider a case where the arrival rate scales with \log SNR. Note that this scaling of arrival process is necessary to ensure a fixed loading and hence a comparable cross-layer interaction as SNR scales.

From the DMT result, we already know that, if the channel operates below the channel ergodic capacity, the asymptotic probability of channel decoding error decays exponentially with SNR. The best one can hope for is that the asymptotic total probability of error decays exponentially with SNR. For that, the asymptotic probability of delay violation needs to decay with SNR. Specifically, we consider a class of independent and identically distributed (i.i.d.)² arrival processes with light tail (i.e., the processes have all moments finite) whose burstiness (defined as the ratio of the standard deviation over the mean of the number of bits arrived at a timeslot) monotonically goes to zero as SNR goes to infinity. We show that for all such processes (called smoothly scaling processes), the total probability of error decays.

The main result of the paper shows that the optimal decay behavior of the asymptotic total probability of bit error depends on how fast the burstiness of the source scales down with SNR. If the source burstiness scales down too slowly (too quickly), the majority of the errors are due to delay violation (channel error), i.e., the total probability of error is asymptotically dominated by delay-violation (channel-error) events. However, at the proper scaling where the burstiness scales linearly with $\frac{1}{\sqrt{\log \text{SNR}}}$ and with the optimal coding duration and transmission rate, the occurrences of channel errors and delay-violation errors are asymptotically balanced. Equivalently, one can interpret our result, the optimal choice of block coding duration and transmission rate, as that which balances the channel atypicality (deep fading or outage events) and the arrival atypicality (large bursts of arrivals).

We apply this result to several examples of outage-limited communication systems to find the optimal setting of the operating parameters.

C. Outline of the Paper

The precise models for the coding/channel process and the bit-arrival/queue process are described in Section II. We precisely define the scaling of the source process with SNR and give a simple example of such source processes. Section III provides

²Note that, since the adopted channel model is not assumed to be i.i.d., assuming an i.i.d. arrival process, intuitively, is not consequential: think of our chosen time slot as an upperbound for the “coherence time” of the arrival process. The i.i.d. source assumption mostly serves to simplify the exposition and presentation of results, and does not fundamentally limit the setting.

the asymptotic probability of delay violation. The main result of the paper is found in Theorem 1 of Section IV. This theorem provides the optimal asymptotic decay rate of the total error probability as well as the optimal coding duration and transmission rate. Section V gives some examples to illustrate the utility of Theorem 1. These examples consider the question of optimally communicating delay sensitive packet stream with a compound Poisson traffic profile over the following outage-limited channels: SISO Rayleigh fast-fading channel, quasi-static cooperative relay channel, and quasi-static MIMO channel. Section VI concludes the paper. Appendices include the proofs of lemmas and theorems.

D. Notations

We use the following symbols and notations. We use ρ to denote SNR. The notation $\stackrel{g}{\approx}$ for a strictly increasing and positive-valued function g represents the equivalence between $y(\rho) \stackrel{g}{\approx} z(\rho)$ and $\lim_{\rho \rightarrow \infty} \frac{\log y(\rho)}{g(\log \rho)} = \lim_{\rho \rightarrow \infty} \frac{\log z(\rho)}{g(\log \rho)}$. We define $\stackrel{g}{\geq}$ and $\stackrel{g}{\leq}$ in a similar manner. Note that when g is an identity function, then $\stackrel{g}{\approx}$ is equivalent to the familiar \doteq notation in the DMT analysis [1].

We denote the high-SNR approximation of the ergodic capacity of AWGN channel by $N := \log \rho$ and use N and $\log \rho$ interchangeably. The sets \mathbb{Z} , \mathbb{N} , and \mathbb{Z}^+ represent the set of all, positive, and non-negative integers, respectively. In addition, the set \mathbb{T} represents the set $\{1, 2, \dots, \lfloor \frac{D}{2} \rfloor\}$. Flooring and ceiling functions are denoted by $\lfloor \cdot \rfloor$ and $\lceil \cdot \rceil$, respectively. For all $a \leq b$, $[x]_a^b = \max(a, \min(b, x))$ and $[x]^+ = \max(x, 0)$. We write $g(x) = \Theta(h(x))$ to denote that the function g scales linearly with the function h , i.e., $\lim_{x \rightarrow \infty} \frac{g(x)}{h(x)} < \infty$ and $\lim_{x \rightarrow \infty} \frac{h(x)}{g(x)} < \infty$. Finally, for any function f , we denote its convex conjugate, f^* , by $f^*(x) = \sup_{\theta \in \mathbb{R}} \theta x - f(\theta)$.

II. SYSTEM MODEL

As discussed in the introduction, we consider a system composed of a bursty and delay-limited information source, concatenated with an infinite buffer and a fading channel, as shown in Fig. 1. We assume the queue follows a first-come-first-serve (FCFS) discipline. The departures out of the queue occur according to a block channel coding scheme, while the arrivals to the queue follow a stochastic model. If the transmission rate is above the instantaneous capacity of the channel, an outage event is said to occur where the received signal is erroneously decoded. The delay requirement asks that each bit of information be decoded at the destination within a maximum allowable delay of D timeslots from the time it arrives at the buffer. Otherwise, the bit will be obsolete, discarded, and counted as erroneous.³ We assume no retransmission of unsuccessful transmissions or those bits which violate the delay bound. In Sections II-A–C, we describe in detail the models for the channel, the arrival process, and the system performance measure.

³Note that due to the constant service rate of the queue and the FCFS service discipline, any bits arriving at the queue know immediately whether they will exceed their delay constraints, using the knowledge of the current queue length. It seems wise to drop these bits immediately after their arrivals to improve the system performance. However, we do not need to consider such method because it has been established (see [14, Theorem 7.10]) that, in the asymptotic regime of interest, such method does not improve the exponent of the delay violation probability.

A. Channel and Coding Model

We consider a general fading-channel model

$$\underline{y} = H\underline{x} + \underline{w}$$

where \underline{x} is the transmitted vector, H is the channel matrix, \underline{y} is the received signal, and \underline{w} is the noise vector. The average SNR is defined as [1]

$$\rho := \frac{\mathbb{E}[\|\underline{H}\underline{x}\|^2]}{\mathbb{E}[\|\underline{w}\|^2]}$$

and in the asymptotic scale of interest, it is equivalent to

$$\rho \doteq \mathbb{E}[\|\underline{x}\|^2].$$

Coding takes place over T timeslots, using rate- R , length- T codes that meet the DMT tradeoff $d_{\text{ch}}(r, T)$ [1], defined as

$$d_{\text{ch}}(r, T) := - \lim_{\rho \rightarrow \infty} \frac{\log P_{\text{ch}}(r, T, \rho)}{\log \rho} \quad (1)$$

where $P_{\text{ch}}(r, T, \rho)$ is the codeword error probability induced by the channel, given an optimal code of *multiplexing gain* r , coding block size T timeslots,⁴ and average SNR ρ . The channel multiplexing gain r is related to the transmission rate R as (refer to [1])

$$r := \lim_{\rho \rightarrow \infty} \frac{R}{\log \rho}. \quad (2)$$

That is, the transmission rate R is assumed to scale linearly as $r \log \rho$. We denote by r_{max} the maximum value of r , i.e., $0 \leq r \leq r_{\text{max}}$. This r_{max} relates to the ergodic capacity as

$$r_{\text{max}} = \mathbb{E}_H \frac{\max_{\underline{x}} I(\underline{x}; \underline{y})}{\log \rho}$$

and is the smallest r such that $d_{\text{ch}}(r, T) = 0$.

The DMT tradeoffs have been extensively studied for various finite-duration communication schemes (for example, see [19], [21]–[24] for MIMO point-to-point communications, [25] for multiple access communications, [26], [27] for cooperative communications, and [20], [28] for cooperative communications with small delay).

Remark 1: The condition that each bit be transmitted over all timeslots in the coding block⁵, together with the first-come first-transmit service discipline, makes it so that every T timeslots, the RT oldest bits are instantaneously removed⁶ from the queue and are transmitted over the next T timeslots. We assume that it is only at the end of the T timeslots that all the RT bits are decoded by the decoder.

⁴For most settings, there exist codes that meet the entire DMT tradeoff in minimum delay, independent of channel dimensionality and fading statistics [19]–[21].

⁵Currently, all known minimum-delay DMT optimal codes over any fading channel with nonzero coefficients ask that each bit be transmitted over each timeslot.

⁶If an insufficient number of bits exists in the buffer, null bits are used and the rate is maintained. It is easy to show that, in the asymptotic scale of interest, the use of null-bits does not incur any change in the SNR exponent of the probability of error.

Example 1 (Rayleigh Fast-Fading SISO Channel): Consider the single-input–single-output (SISO) time-selective channel with Rayleigh fading coefficients (correlated or uncorrelated) and with additive white Gaussian noise at the receiver. The corresponding channel model over T timeslots is given by

$$\underline{y} = \text{diag}(\underline{h}) \underline{x} + \underline{w}$$

where \underline{y} , \underline{h} , \underline{x} , and \underline{w} are $T \times 1$ vectors and $H = \text{diag}(\underline{h})$ is a $T \times T$ diagonal fading matrix with the fading at the t th timeslot, h_t , as its (t, t) element. The optimal DMT, given optimal signaling, takes the form

$$\begin{aligned} d_{\text{ch}}(r, T) &:= - \lim_{\rho \rightarrow \infty} \frac{\log \Pr(I(\underline{x}; \underline{y} | \underline{h}) < 2^{RT})}{\log \rho} \\ &= - \lim_{\rho \rightarrow \infty} \frac{\log \Pr(\prod_{t=1}^T (1 + \rho |h_t|^2) < \rho^{rT})}{\log \rho}. \end{aligned}$$

For the fast-fading case where the coherence time is equal to one timeslot and the elements of \underline{h} are Rayleigh i.i.d. random variables, the tradeoff takes the form

$$d_{\text{ch}}(r, T) = T(1 - r) \quad (3)$$

and it can be met entirely in T timeslots (see [1]). This SISO channel allows for $r_{\text{max}} = 1$.

Other examples which will be discussed later in Section V are quasi-static MIMO and cooperative-relay channels. In this paper, for simplicity we assume that $d_{\text{ch}}(r, T)$ is continuous on r , decreasing on r , and increasing on T .

B. Smoothly Scaling Bit-Arrival Process

In this section, we describe the SNR-scaling of a family of arrival processes of interest. The specific choice of SNR-scaling for the statistics of the bit-arrival process is such that the average traffic load of the system (defined as the ratio of the average arrival rate over the ergodic capacity) is kept constant, independent of SNR.⁷ This means that scaling in the ergodic capacity $r_{\text{max}} \log \rho$ ($= r_{\text{max}} N$) is matched by scaling the average bit-arrival rate as $\lambda \log \rho$ ($= \lambda N$) as well, for some $\lambda > 0$. Now we are ready to introduce the arrival process of interest: The sequence of asymptotically *smoothly scaling* bit-arrival processes, in which the process becomes “smoother” for increasing N .

Definition 1: Let \mathcal{G} denote a class of functions which contains any function $g : \mathbb{R}^+ \mapsto \mathbb{R}^+$ (called *scaling function*) which is continuous and strictly increasing and whose tail behavior is such that

$$\lim_{x \rightarrow \infty} \frac{g(x)}{\log x} = \infty. \quad (4)$$

⁷It can be seen that unless the traffic load (average bit arrival rate over the channel rate) scales as $\log(\text{SNR})$, i.e., $\lim_{\rho \rightarrow \infty} \frac{\mathbb{E}[A_t]}{\log \rho} = \ell$ for some fixed $0 < \ell < \infty$, the problem is void of cross-layer interactions. Otherwise if $\ell = 0$, corresponds to the case where too few bits arrive and effectively there is no queuing delay. On the other hand, when the traffic load scales much faster than $\log(\text{SNR})$, i.e., $\ell = \infty$, the overall performance is dominated by queuing delay, independently of the channel characteristics.

Definition 2: (g -Smoothly Scaling Source): Consider a scaling function $g \in \mathcal{G}$ and a family of bit-arrival processes $(A^{(N)}, N \in \mathbb{N})$, where $A^{(N)} = (A_t^{(N)}, t \in \mathbb{Z})$ denotes an i.i.d. sequence of the random numbers $A_t^{(N)}$ of bits that arrive at time t with $E[A_t^{(N)}] = \lambda N$, for all t . The family of bit-arrival processes is said to be g -smoothly scaling if the limiting g -scaled logarithmic moment generating function, defined for each $\theta \in \mathbb{R}$ as

$$\Lambda(\theta) = \lim_{N \rightarrow \infty} \frac{\log E[\exp(\frac{\theta g(N)}{N} A_1^{(N)})]}{g(N)} \quad (5)$$

exists as an extended real number in $\mathbb{R}^* := \mathbb{R} \cup \{\infty\}$ and is finite in a neighborhood of the origin, essentially smooth, and lower-semicontinuous [14].⁸

Remark 2: It is straight forward to show that Λ is convex and $\Lambda'(0) = \lambda$ (see [14, Lemma 1.11]).

Note that λ describes how close the average bit-arrival rate is to the asymptotic approximation of the ergodic capacity of the channel. For stability purpose and to ensure the existence of a stationary distribution, we require that $\lambda < r_{\max}$. Also, note that we abuse the notation and denote the arrival process by $A_t^{(N)}$, despite its possible dependency on the scaling function g .

Motivation for Smoothly Scaling Assumption: The assumption of g -smoothly scaling arrival processes allows us to find the decay rate of the tail probability of the sequence of process $(S_t^{(N)}, N \in \mathbb{N})$, which is a sum process defined as

$$S_t^{(N)} = \sum_{j=1}^t A_j^{(N)}, \quad t \in \mathbb{N}$$

since $(A_j^{(N)}, j \in \mathbb{Z})$ are i.i.d., $S_t^{(N)}$ is also a g -smoothly scaling process with the limiting g -scaled log moment generating function Λ_{S_t} given as

$$\Lambda_{S_t}(\theta) := \lim_{N \rightarrow \infty} \frac{\log E[\exp(\frac{\theta g(N)}{N} S_t^{(N)})]}{g(N)} = t\Lambda(\theta). \quad (6)$$

Now, given that the sequence $(S_t^{(N)}, N \in \mathbb{N})$ is g -smoothly scaling, we can use the Gärtner–Ellis theorem (see, e.g., [10] and [14]) to give the following result on the decay rate of the tail probability of the sequence. The following proposition provides an important basis for the analysis of the asymptotic probability of delay violation in Section III.

Proposition 1: (Gärtner–Ellis Theorem for g -Smoothly Scaling Process): Consider $g \in \mathcal{G}$ and a family of g -smoothly scaling processes $(A^{(N)}, N \in \mathbb{N})$ with the limiting g -scaled log moment generation function Λ . Let $S_t^{(N)} = \sum_{i=1}^t A_i^{(N)}$, for $t \in \mathbb{N}$. Then, for $a > \lambda t$, we have

$$\lim_{N \rightarrow \infty} \frac{1}{g(N)} \log P\left(\frac{S_t^{(N)}}{N} > a\right) = -t\Lambda^*(a/t) \quad (7)$$

where Λ^* is the convex conjugate of Λ .

Proof: See Appendix A. \square

⁸A function $f : \mathbb{R} \mapsto \mathbb{R}^*$ is essentially smooth if the interior of its effective domain $\mathcal{D} = \{x : f(x) < \infty\}$ is nonempty, if it is differentiable in the interior of \mathcal{D} and if f is steep, which means that for any sequence θ_n which converges to a boundary point of \mathcal{D} , then $\lim_{n \rightarrow \infty} |f'(\theta_n)| = +\infty$. f is lower semicontinuous if its level sets $\{x : f(x) \leq \alpha\}$ are closed for $\alpha \in \mathbb{R}$.

1) Asymptotic Characteristic of Smoothly Scaling Processes: Intuitively, the g -smoothly scaling arrival processes become smoother as SNR increases. This intuition follows from (5), which implies that for $\theta \in \mathbb{R}$ such that $\Lambda(\theta) < \infty$ and $\epsilon > 0$, there exists N_0 such that for $N > N_0$

$$\begin{aligned} \exp(g(N)\Lambda(\theta) - g(N)\epsilon) &< E\left[\exp\left(\frac{\theta g(N)}{N} A_1^{(N)}\right)\right] \\ &< \exp(g(N)\Lambda(\theta) + g(N)\epsilon). \end{aligned}$$

Then, if we let $Y_{g(N)}$ be a sum of $g(N)$ i.i.d. random variables (i.e., $Y_{g(N)} := X_1 + \dots + X_{g(N)}$ with $E[e^{\theta X_1}] = e^{\Lambda(\theta)}$), we have $E[e^{\theta(Y_{g(N)})}] = e^{\Lambda(\theta)g(N)}$. Therefore, at sufficiently large N , $\frac{g(N)}{N} A_t^{(N)}$ and $Y_{g(N)}$ have the same moment generating function and hence the same distribution. If we define the burstiness of the random variable $A_1^{(N)}$ as the (dimensionless) ratio of its standard deviation over its mean,⁹ then, using the above intuition, the burstiness $\frac{\text{std}(A_1^{(N)})}{E[A_1^{(N)})]}$ for large N is approximately equal to $\frac{\text{std}(\frac{g(N)}{N} \sum_{i=1}^{g(N)} X_i)}{\lambda N}$, which is reduced to $\frac{\text{std}(X_1)}{\lambda \sqrt{g(N)}}$. Hence, the burstiness of $A_1^{(N)}$ decays to zero approximately as $\frac{1}{\sqrt{g(\log \rho)}}$. In other words, the g -smoothly scaling arrival processes become smoother as SNR increases.

Examples of Smoothly Scaling Processes: One of the common arrival processes used for traffic modeling is a compound Poisson process with exponential packet size, denoted as CPE. For this source, the random number of bits, $A_t^{(N)}$, arrived at timeslot t , is i.i.d. across time t and is in the form of

$$A_t^{(N)} = \sum_{i=1}^{M_t^{(N)}} Y_{i,t}^{(N)} \quad (8)$$

where $M_t^{(N)}$ is the random variable corresponding to the number of packets that have arrived at the t th timeslot, and where $Y_{i,t}^{(N)}$ corresponds to the random number of bits in the i th packet. $M_t^{(N)}$ are independently drawn from a Poisson distribution with mean $\nu(N)$; and $Y_{i,t}^{(N)}$, $i = 1, \dots, M_t^{(N)}$, are independently drawn from an exponential distribution with mean $\frac{1}{\mu(N)}$ (nats per packet). Note that the assumption that $E[A_t^{(N)}] = \lambda N$ forces that $\frac{\nu(N)}{\mu(N)} = \lambda N$. In addition, a larger average packet size $\frac{1}{\mu(N)}$ implies a more bursty arrival process.¹⁰ It is known (see [16]) that the log moment generating function of this CPE random variable $A_t^{(N)}$ is

$$\log E[e^{\theta A_t^{(N)}}] = \begin{cases} \frac{\theta \nu(N)}{\mu(N) - \theta}, & \theta < \mu(N) \\ \infty, & \text{otherwise.} \end{cases} \quad (9)$$

The following examples illustrate that, depending on the scaling of the average packet arrival rate and the average packet

⁹Note that the burstiness definition here is basically the normalized variation of the random variable around its typical value (its mean). A more familiar definition of traffic burstiness would involve how the traffic are correlated with time, i.e., a bursty source tends to have large bursts of arrivals in a short period of time. However, since we only consider the source which is i.i.d. over time, we use this definition of burstiness.

¹⁰It can be easily shown that the burstiness of this CPE process, as defined in Section II-B2, is $\sqrt{\frac{2}{\lambda N \mu(N)}}$.

size, some CPE processes may or may not be g -smoothly scaling.

Example 2 (g -Smoothly Scaling CPE Process): For $g \in \mathcal{G}$ and $\mu > 0$, consider a CPE process $A_t^{(N)}$ with packet arrival rate $\mu\lambda g(N)$ and average packet size $\frac{N}{\mu g(N)}$. This family of processes is g -smoothly scaling because, using (9), we have

$$\Lambda(\theta) := \lim_{N \rightarrow \infty} \frac{\log E[e^{\frac{\theta g(N)}{N} A_t^{(N)}}]}{g(N)} = \begin{cases} \frac{\mu\lambda\theta}{\mu-\theta}, & \theta < \mu \\ \infty, & \text{otherwise} \end{cases} \quad (10)$$

which satisfies the conditions in the definition of g -smoothly scaling. Since we will use this particular g -smoothly scaling CPE process for examples in the paper, we denote it as $\text{CPE}(\lambda, \mu, g, N)$. It is useful to note a particular case when $g(N)$ grows linearly with N . Using a property of the Poisson process [29], this particular scaling case can be viewed as aggregating an increasing number of Poisson traffic streams (this number grows linearly with N), with each stream having the same packet length distribution.

To complete our discussion on smoothly scaling processes, below we give an example of a family of CPE arrival processes which is not g -smoothly scaling.

Example 3: A family of CPE processes where $A_t^{(N)}$ has packet arrival rate $\mu\lambda$ and average packet size N/μ (note the dependence on N only in the average packet size) is not g -smoothly scaling for any $g \in \mathcal{G}$. This is because, using (9), we have

$$\lim_{N \rightarrow \infty} \frac{\log E[e^{\frac{\theta g(N)}{N} A_t^{(N)}}]}{g(N)} = \begin{cases} 0, & \theta \leq 0 \\ \infty, & \text{otherwise} \end{cases}$$

which is not finite in the (open) neighborhood of $\theta = 0$. Hence, this family of processes is not g -smoothly scaling.

Remark 3: The scaling function, g , describes the way the source statistics scale with SNR. Example 2 describes the case of the compound Poisson process, where g can be identified as the function that specifies how the average packet arrival rate ($\mu\lambda g(\log \text{SNR})$) and the average packet size ($\frac{\log \text{SNR}}{\mu g(\log \text{SNR})}$) scale with SNR.

C. Performance Measure and System Objective

The overall performance measure is the total probability of bit loss, $P_{\text{tot}}(r, T)$, where loss can occur due to channel decoding error or the end-to-end delay violation. Specifically,

$$P_{\text{tot}}(r, T) := P_{\text{ch}}(r, T) + (1 - P_{\text{ch}}(r, T)) P_{\text{delay}}(r, T) \quad (11)$$

where $P_{\text{ch}}(r, T)$ denotes the probability of decoding error due to channel outage and $P_{\text{delay}}(r, T)$ denotes the probability of delay violation. We are interested in finding the high-SNR asymptotic approximation of $P_{\text{tot}}(r, T)$ as a function of r, T, SNR, D , as well as the source and channel statistics (including λ and the source scaling function g). In the interest of brevity, we denote P_{tot} as a function of only r and T , the two parameters over which the performance will later be optimized.

Since the high-SNR asymptotic expression of $P_{\text{ch}}(r, T)$ is already given by the DMT in (1), what remains is to find the

asymptotic expression for $P_{\text{delay}}(r, T)$, which is shown in Section III.

III. ASYMPTOTIC ANALYSIS OF PROBABILITY OF DELAY VIOLATION

In this section, we derive the asymptotic probability of delay violation $P_{\text{delay}}(r, T)$ for the channel multiplexing rate r and coding block size T . We observe that the adopted block coding forces the queue to have a *batch service* that occurs every T timeslots with the instantaneous removal of the oldest rNT bits. The decay rate of the asymptotic tail probability of the sum arrival process, given in Proposition 1, in conjunction with an asymptotic analysis of a queue with deterministic batch service, gives the following result.

Lemma 2: Given $g \in \mathcal{G}, T \in \mathbb{T}, r > \lambda$, a batch service of rNT every T timeslots, and a g -smoothly scaling bit-arrival process characterized by the limiting g -scaled log moment generation function Λ , the decay rate of $P_{\text{delay}}(r, T)$ is given by the function I , i.e.,

$$\lim_{N \rightarrow \infty} \frac{1}{g(N)} \log P_{\text{delay}}(r, T) = -I(r, T), \quad (12)$$

where

$$I(r, T) = \min_{\substack{t \in \mathbb{Z}^+ \\ tT + T - 1 - k > 0}} (tT + T - 1 - k) \times \Lambda^* \left(r + \frac{(D + 1 - 2T)r}{tT + T - 1 - k} \right) \quad (13)$$

for $k = D(\text{mod } T)$. In addition, $I(r, T)$ is lower-semicontinuous and increasing on r .

Proof: See Appendix B. \square

Approximation 1: Relaxing the integer constraint in (13) gives the lower bound of I as

$$I(r, T) \geq \delta_r r (D + 1 - 2T) =: I_{ir}(r, T) \quad (14)$$

where

$$\delta_r = \sup\{\theta > 0 : \Lambda(\theta) < \theta r\}. \quad (15)$$

We use this lower bound as an approximation to I as well, i.e.,

$$I(r, T) \approx I_{ir}(r, T) = \delta_r r (D + 1 - 2T). \quad (16)$$

Proof: See Appendix B. \square

Example 4: For a g -smoothly scaling $\text{CPE}(\lambda, \mu, g, N)$ bit-arrival process, the function I in (13) can be calculated exactly with the following Λ^* :

$$\Lambda^*(x) = \mu \left(\sqrt{x} - \sqrt{\lambda} \right)^2, \quad x \in \mathbb{R}. \quad (17)$$

However, an approximation of I in (16) is simpler to work with and given as

$$I(r, T) \approx I_{ir}(r, T) = \mu(r - \lambda)(D + 1 - 2T), \quad (18)$$

where, using (15) and (10), δ_r is given as

$$\delta_r = \mu \left(1 - \frac{\lambda}{r} \right). \quad (19)$$

We will see via numerical examples in Section V-A1 that the approximation in (18) is sufficient for our purpose.

IV. MAIN RESULT: OPTIMAL ASYMPTOTIC TOTAL PROBABILITY OF ERROR

In this section, we present the main result of the paper which states the optimal decay rate of the high-SNR asymptotic total probability of bit error. Recall the definition of P_{tot} from (11):

$$P_{\text{tot}}(r, T) := P_{\text{ch}}(r, T) + (1 - P_{\text{ch}}(r, T))P_{\text{delay}}(r, T)$$

where we now know that

$$P_{\text{ch}}(r, T) \doteq \rho^{-d_{\text{ch}}(r, T)}$$

and

$$P_{\text{delay}}(r, T) \stackrel{g}{\doteq} e^{-I(r, T)g(\log \rho)}.$$

Hence, the asymptotic optimal decay behavior of P_{tot} depends on the function g . The following theorem gives the main result of the paper.

Theorem 1: Consider $g \in \mathcal{G}$ and a g -smoothly scaling bit-arrival process. The optimal rate of decay of the asymptotic probability of total bit error, maximized over all $r \in (\lambda, r_{\text{max}})$ and $T \in \mathbb{T}$, and the optimizing r^* and T^* are given, depending on the tail behavior of the function g , as follows.

Case 1: If $\lim_{N \rightarrow \infty} \frac{g(N)}{N} = \gamma \in (0, \infty)$, then

$$\begin{aligned} d^* &:= \sup_{\substack{r \in (\lambda, r_{\text{max}}) \\ T \in \mathbb{T}}} \lim_{\rho \rightarrow \infty} \frac{-\log P_{\text{tot}}(r, T)}{\log \rho} \\ &= d_{\text{ch}}(r^*, T^*) = \gamma I(r^*, T^*) \end{aligned} \quad (20)$$

where

$$r^*(T) := \inf\{r \in (\lambda, r_{\text{max}}) : \gamma I(r, T) = d_{\text{ch}}(r, T)\} \quad (21)$$

$$T^* = \arg \max_{T \in \mathbb{T}} I(r^*(T), T) \quad (22)$$

$$r^* = r^*(T^*). \quad (23)$$

Case 2: If $\lim_{N \rightarrow \infty} \frac{g(N)}{N} = 0$ and $\lim_{N \rightarrow \infty} \frac{g(N)}{\log N} = \infty$, then

$$\sup_{r \in (\lambda, r_{\text{max}}), T \in \mathbb{T}} \lim_{\rho \rightarrow \infty} \frac{-\log P_{\text{tot}}(r, T)}{g(\log \rho)} \leq \max_{T \in \mathbb{T}} I(r_{\text{max}}, T). \quad (24)$$

Case 3: If $\lim_{N \rightarrow \infty} \frac{g(N)}{N} = \infty$, then

$$\sup_{r \in (\lambda, r_{\text{max}}), T \in \mathbb{T}} \lim_{\rho \rightarrow \infty} \frac{-\log P_{\text{tot}}(r, T)}{\log \rho} \leq d_{\text{ch}}\left(\lambda, \left\lfloor \frac{D}{2} \right\rfloor\right). \quad (25)$$

Proof: See Appendix C. \square

Theorem 1 shows that the optimal decay behavior of the asymptotic total probability of error depends on the tail behavior of the function g . As discussed earlier, the burstiness of the g -smoothly scaling arrival process scales down as $\Theta\left(\frac{1}{\sqrt{g(\log \rho)}}\right)$. Below, we discuss each case of Theorem 1, with respect to the scaling of the source burstiness:

In Case 1, where the source burstiness scales down with $\Theta\left(\frac{1}{\sqrt{\log \rho}}\right)$, both components of the probability of error decay exponentially with $\log \rho$. In this setting, one can optimize the choices of r and T to arrive at a nontrivial optimal decay rate d^* . The optimal r^* and T^* balance and minimize the decay rate in $P_{\text{ch}}(r, T)$ and $P_{\text{delay}}(r, T)$. Hence, for Case 1, the optimal asymptotic total probability of error decays as follows:

$$P_{\text{tot}}(r^*, T^*) \doteq P_{\text{delay}}(r^*, T^*) \doteq P_{\text{ch}}(r^*, T^*) \doteq \rho^{-d^*}.$$

Note that d^* is nothing but the optimal negative SNR exponent.

In Case 2, where the source burstiness scales down slower than $\Theta\left(\frac{1}{\sqrt{\log \rho}}\right)$ but faster than $\Theta\left(\frac{1}{\sqrt{\log \log \rho}}\right)$, we have that $P_{\text{tot}}(r, T)$ is asymptotically equal to $P_{\text{delay}}(r, T)$ for all $r \in (\lambda, r_{\text{max}})$ and $T \in \mathbb{T}$. In this case, the decay rate of $P_{\text{tot}}(r, T)$ is equal to $I(r, T)$. In other words, the channel error (outage) probability is dominated by the delay violation probability and, hence, can be ignored.

Finally, in Case 3, when the source burstiness scales down faster than $\Theta\left(\frac{1}{\sqrt{\log \rho}}\right)$, we have the opposite of Case 2. In Case 3, the delay violation probability is dominated by the channel error probability and, hence, can be ignored.

A. Approximation of the Optimal Negative SNR Exponent

For Case 1 in Theorem 1, we use the following approximation which is an immediate result of relaxing the integer-constrained optimizations of I and T^* to obtain approximated expressions with much simpler forms. These approximations become especially useful in Section V.

Approximation 2: Relaxing the integer constraints in the calculation of I (as in Approximation 1) and T^* in (22) gives the following ‘‘integer-relaxed’’ approximations for d^* , r^* , and T^* :

$$\begin{aligned} d^* &\approx d_{ir}^* := d_{\text{ch}}(r_{ir}^*, T_{ir}^*) \\ T^* &\approx T_{ir}^*, \text{ and } r^* \approx r_{ir}^*, \end{aligned} \quad (26)$$

where, for δ_r given in (15) and any $T \in \mathbb{T}$,

$$r_{ir}^*(T) := \min\{r \in (\lambda, r_{\text{max}}) : d_{\text{ch}}(r, T) = \gamma \delta_r (D - 2T + 1)\} \quad (27)$$

and

$$T_{ir}^* = \left\lceil \min \left\{ T \in \mathbb{R}^+ : \frac{d}{dT} (d_{\text{ch}}(r_{ir}^*(T), T)) = 0 \right\} \right\rceil^{\lfloor \frac{D}{2} \rfloor} \quad (28)$$

$$r_{ir}^* = r_{ir}^*(T_{ir}^*). \quad (29)$$

V. APPLICATIONS OF THE RESULT

In this section, we apply the result of Case 1 in Theorem 1 to analyze and optimize the end-to-end error probability of systems communicating delay-sensitive and bursty traffic over three outage-limited channels: SISO Rayleigh fast-fading

channel, quasi-static cooperative relay channel, and quasi-static MIMO channel.

To illustrate the methodology, we restrict our attention to the case of CPE($\lambda, \mu, g, \log \rho$) arrival process where $g(\log \rho) = \log \rho$, for simplicity. Note that to better gain insights, we use the integer-relaxed approximations obtained in Approximation 2.

A. SISO Rayleigh Fast-Fading Channel

Our first example considers an example of SISO Rayleigh fast-fading channel, whose $d_{\text{ch}}(r, T) = T(1-r)$ (see (3)). Combining this with (19) and (27) gives the optimal choice of multiplexing gain when the coding duration is fixed at T as

$$r_{ir}^*(T) = \lambda + \frac{1-\lambda}{1 + \frac{\mu(D+1-2T)}{T}}. \quad (30)$$

In addition, using (28), the integer-relaxed approximated optimal coding duration can be expressed as

$$T_{ir}^* = \left\lceil \frac{1}{1 + \frac{1}{\sqrt{2\mu}}} \frac{D+1}{2} \right\rceil_{1}^{\lfloor \frac{D}{2} \rfloor}. \quad (31)$$

Inserting T_{ir}^* into (30), we get the approximated optimal channel multiplexing gain as

$$r_{ir}^* = r_{ir}^*(T_{ir}^*) = \left[\lambda + \frac{1-\lambda}{1 + \sqrt{2\mu}} \right]_{r_{ir}^*(1)}^{r_{ir}^*(\lfloor \frac{D}{2} \rfloor)}. \quad (32)$$

Also, from (26), the approximated optimal negative SNR exponent is given as

$$\begin{aligned} d_{ir}^* &= T_{ir}^* (1 - r_{ir}^*) \\ &= \left[\frac{1}{\left(1 + \frac{1}{\sqrt{2\mu}}\right)^2} \frac{(D+1)}{2} (1-\lambda) \right]_{1-r_{ir}^*(1)}^{\lfloor \frac{D}{2} \rfloor (1-r_{ir}^*(\lfloor \frac{D}{2} \rfloor))}. \end{aligned} \quad (33)$$

Below, we provide some observations of the above results:

- The above result on d^* can also be interpreted as a tradeoff which describes the relation between the normalized average arrival rate

$$\lambda := \lim_{N \rightarrow \infty} (\text{average bit-arrival rate})/N = \lim_{N \rightarrow \infty} \frac{\mathbb{E}[A_t^N]}{N}$$

and the corresponding optimal negative SNR exponent $d_{ir}^*(\lambda)$ as a function of the delay bound D , and the average packet size $1/\mu$. For constant bit arrivals (CBR) at rate $\lambda \log \rho$, i.e., mathematically when $1/\mu \rightarrow 0$, any coding durations less than half¹¹ of D (or more precisely $\lfloor \frac{D}{2} \rfloor$) and any channel multiplexing rates greater than λ result in zero probability of delay violation. Hence, the optimal negative SNR exponent of the total error probability, denoted by d_{CBR}^* , is equal to the corresponding channel

¹¹The first half of D is spent waiting for the next coding block and the other half waiting to be decoded at the end of the block.

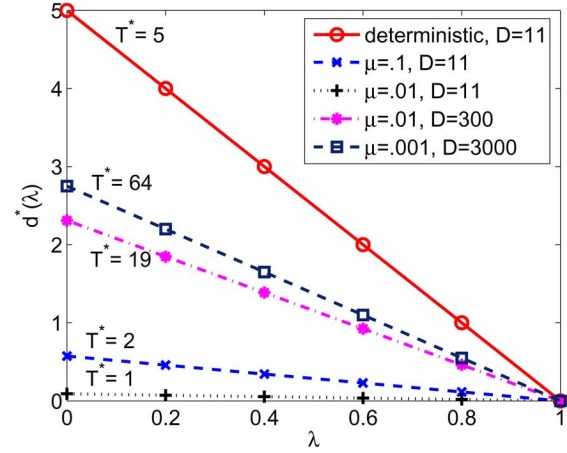


Fig. 2. SISO, Rayleigh fast-fading, coherent channel. The solid line describes the DMT ($r = \lambda$). The dashed and dotted lines describe $d_{ir}^*(\lambda)$ for various μ and D .

diversity when the optimal coding duration is at its maximum value, $\lfloor \frac{D}{2} \rfloor$, and the channel multiplexing gain is at its minimum, λ . That is

$$d_{CBR}^*(\lambda) = \left\lfloor \frac{D}{2} \right\rfloor (1 - \lambda).$$

It is not surprising that this coincides with the classical DMT. With traffic burstiness, however, the optimal negative SNR exponent $d_{ir}^*(\lambda)$ given in (33) is smaller than $d_{CBR}^*(\lambda)$. The ratio

$$\frac{d_{ir}^*(\lambda)}{d_{CBR}^*(\lambda)} \approx \frac{1}{\left(1 + \frac{1}{\sqrt{2\mu}}\right)^2} \leq 1$$

can be interpreted as the reduction factor on the SNR exponent in the presence of burstiness. Fig. 2 shows the impact of traffic burstiness (which is parameterized by μ) on $d_{ir}^*(\lambda)$.

- From a coding point of view, T_{ir}^* is independent of the average bit-arrival rate λ . This implies that for a fixed value of the average packet size $1/\mu$, the optimal negative SNR exponent is achieved by a fixed-duration $1 \times T_{ir}^*$ code. Optimal codes for this setting exist for all values of r and T ([20], [21], [28]). On the other hand, if T is already given, the performance is optimized when the coding multiplexing gain is chosen as in (30), i.e.,

$$r_{ir}^*(T) = \lambda + \frac{1-\lambda}{1 + \frac{\mu(D+1-2T)}{T}}.$$

- Since $r_{\text{max}} = 1$ for this SISO channel, we can verify that $r_{ir}^* \nearrow r_{\text{max}}$ for very bursty traffic (i.e., $1/\mu \rightarrow \infty$). That is for very bursty traffic the channel should operate close to its highest possible rate, which is the channel ergodic capacity.

1) *Numerical Comparison of the Approximation:* Before we move to the next example, we illustrate numerically that the approximations in (31)–(33) well approximate their actual values in Theorem 1. In Fig. 3, we show an example of a comparison at $1/\mu = 100$ and various values of D and λ . We observe that

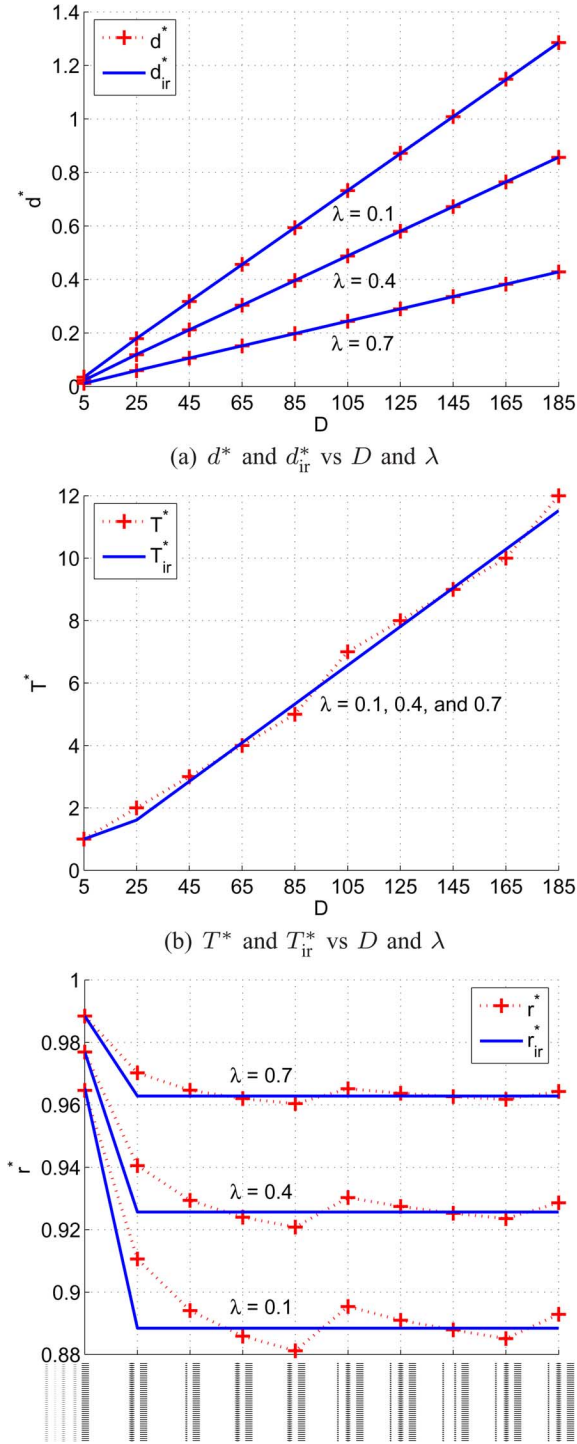


Fig. 3. Comparisons of the exact values d^* , T^* , and r^* (3a, 3b, 3c, respectively) and the integer-relaxed approximations d_{ir}^* , T_{ir}^* , and r_{ir}^* , at various D and λ . The dotted lines with markers correspond to the exact solutions while the solid lines represent the approximated solutions.

the approximated values match well with the exact values if D is sufficiently large. The matching is remarkably good for d^* and d_{ir}^* . Note that r_{ir}^* is independent of D , except when D is so small that $T_{ir}^* = 1$.

B. Cooperative Wireless Networking With Optimal Clustering

As studied in [30], [31], we consider communicating bursty and delay-limited information from an information source in a

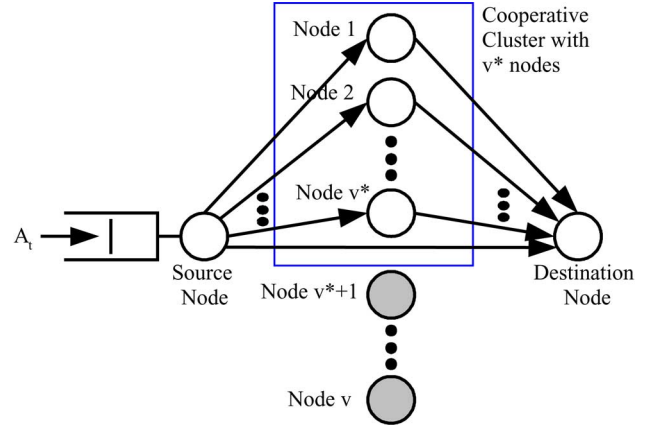


Fig. 4. Snapshot of a wireless network, where the source node utilizes a subset of its peers (nodes 1, 2, ..., v^*) as relays for communicating with the destination.

cooperative wireless relay network, shown in Fig. 4, where the diversity benefit of user cooperation is due to encoding across space and time [26], [32]. In the absence of delay limitation, having more cooperative users almost always improves performance. This is not the case, though, when one considers burstiness and delay QoS requirement. Take for example a network where the information-source node cooperates with v relays, under an orthogonal amplify-and-forward (OAF) cooperative diversity scheme and half-duplex constraint. This cooperation scheme gives the DMT:

$$d_{ch}^{coop}(r) = (v + 1)(1 - 2r).$$

Note that $r_{max} = 1/2$ under this protocol. To realize this amount of diversity, the coding duration T is required to be at least $2(v + 1)$ channel uses or timeslots. This means that, in spite of the increase in the negative SNR exponent of the probability of decoding error with the number of cooperative relays, relaying over all nodes in the network might not be desirable as it increases the delay violations. Applying the result of Approximation 2 to CPE source and the above $d_{ch}^{coop}(r)$ with $T = 2(v + 1)$, the optimal performance is achieved when the nodes cooperate in clusters with

$$v^* \approx v_{ir}^* = \left\lceil \frac{D + 1}{4 \left(1 + \frac{1}{\sqrt{2\mu}}\right)} - 1 \right\rceil$$

relays and transmit at multiplexing rate,

$$r^* \approx r_{ir}^* = \frac{1}{2} - \frac{\frac{1}{2} - \lambda}{1 + \frac{1}{\sqrt{2\mu}}}.$$

Note that v_{ir}^* is independent of the traffic average arrival rate λ . This means that meeting the optimal tradeoff for various values of λ does not require modifying the cluster sizes, unless the traffic burstiness (parameterized by the average packet size $1/\mu$) changes.

C. MIMO Quasi-Static Communications

In the case of the MIMO Rayleigh-fading channel with n_t transmit and n_r receive antennas, and with complete channel

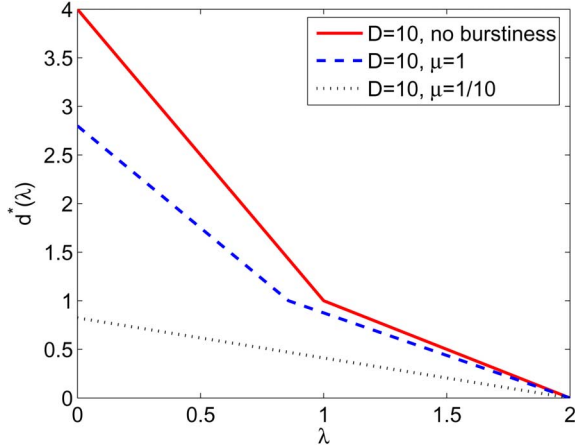


Fig. 5. MIMO, quasi-static, coherent 2×2 channel. d^* versus λ for different values of D and μ .

state information at the receiver (CSIR) and no CSI at the transmitter, the channel diversity gain $d_{\text{ch}}(r)$ is shown (see [1]) to be a piecewise linear function that connects points

$$(k, (n_t - k)(n_r - k)), \quad k = 0, 1, \dots, \min(n_t, n_r). \quad (34)$$

The entire tradeoff is met if $T \geq n_t$ [19]. An example of the effect of burstiness is shown in Fig. 5, for the case of the 2×2 Rayleigh-fading channel ($n_t = n_r = 2$). By assuming that T is given (not an optimizing parameter) and equal to 2, the optimal multiplexing gain r^* , which balances the SNR exponents of the probabilities of delay violation and decoding error, is the solution to $d_{\text{ch}}(r^*) = I(r^*, T = 2)$. Using the approximation (18) of I for CPE source, the approximation r_{ir}^* is the solution to

$$d_{\text{ch}}(r_{ir}^*) - \mu(r_{ir}^* - \lambda)(D - 3) = 0$$

where d_{ch} is the piecewise linear function connecting points in (34). In other words, r_{ir}^* is given as

$$r_{ir}^* = \begin{cases} \lambda + \frac{2-\lambda}{1+\mu(D-3)}, & \text{if } \lambda \in [1 - \frac{1}{\mu(D-3)}, 2), \\ \lambda + \frac{4-3\lambda}{3+\mu(D-3)}, & \text{if } \lambda \in (0, 1 - \frac{1}{\mu(D-3)}]. \end{cases}$$

Fig. 5 shows the resulting $d^*(\lambda) = d_{\text{ch}}(r_{ir}^*)$ for various values of burstiness μ and D .

VI. SUMMARY AND FUTURE WORK

This work offers a high-SNR asymptotic error performance analysis for communications of delay-limited and bursty information over an outage-limited channel, where errors occur either due to delay or due to erroneous decoding. The analysis focuses on the case where there is no CSIT and no feedback, and on the static case of fixed rate and fixed length of coding blocks. This joint queue-channel analysis is performed in the asymptotic regime of high-SNR and in the assumption of smoothly scaling (with SNR) bit-arrival processes. The analysis provides closed-form expressions for the error performance, as a function of the channel and source statistics. These expressions identify the scaling regime of the source and channel statistics in which either delay or decoding errors are the dominant cause of errors, and the scaling regime in which a prudent choice of the coding duration and rate manages to balance and minimize these

errors. That is, in this latter regime, such optimal choice manages to balance the effect of channel atypicality and burstiness atypicality. To illustrate the results, we provide different examples that apply the results in different communication settings. We emphasize that the results hold for any coding duration and delay bound.

Many interesting extensions of the current work remain. One example is the high-SNR analysis of systems with retransmission mechanism and/or adaptive adjustment of the transmission rate and coding duration as a function of the current queue length at the transmitter. With retransmission, the diversity of the channel can be improved considerably [33] but at the cost of longer and random transmission delays. On the other hand, we may be able to improve the system performance by adjusting the transmission rate according to the need of the bits in the queue. For example, when the queue length is short, we may reduce the transmission rate, which improves the probability of channel decoding error but possibly at the cost of longer delays of the bits that arrive later. However, since in high-SNR analysis the probability of error is asymptotically dominated by the worst case probability, it is not clear whether such adaptive transmission rate mechanism will improve the asymptotic decay rate of the probability of bit error.

In addition, this work focuses on the notion of SNR error exponent as a measure of performance. This view of communication systems provides a tractable and intuitive characterization of various suggested schemes in the high-SNR regime. It would also be interesting to fine-tune the high-SNR asymptotic analysis presented here, for the regime of finite SNR, as well as extend it to different families of bit-arrival processes.

APPENDIX A

PROOF OF PROPOSITION 1

Proposition 1: Consider a g -smoothly scaling process $A^{(N)}$ with the limiting g -scaled log moment generation function Λ . Let $S_t^{(N)} = \sum_{i=1}^t A_i^{(N)}$, for $t \in \mathbb{N}$. Then, for $a > \lambda t$, we have

$$\lim_{N \rightarrow \infty} \frac{1}{g(N)} \log P \left(\frac{S_t^{(N)}}{N} > a \right) = -t\Lambda^*(a/t) \quad (35)$$

where Λ^* is the convex conjugate of Λ .

Proof: Let $n = g(N)$ and $Y_t^{(n)} = \frac{g(N)}{N} S_t^{(N)}$. From (6) and the property of Λ for the g -smoothly scaling process, we have

$$\Lambda_{Y_t}(\theta) := \lim_{n \rightarrow \infty} \frac{1}{n} \log E[e^{\theta Y_t^{(n)}}] = \Lambda_{S_t}(\theta) = t\Lambda(\theta)$$

which exists for each $\theta \in \mathbb{R}$ as an extended real number and is finite in a neighborhood of $\theta = 0$, essentially smooth, and lower-semicontinuous. Then, the Gärtner-Ellis theorem ([14, Theorem 2.11]) shows that $Y_t^{(n)}/n$ (which, in this case, is equivalent to $S_t^{(N)}/N$) satisfies the large deviations principle (LDP) in \mathbb{R} with good convex rate function

$$\Lambda_{Y_t}^*(x) := \sup_{\theta \in \mathbb{R}} \theta x - \Lambda_{Y_t}(\theta) = \sup_{\theta \in \mathbb{R}} \theta x - t\Lambda(\theta) = t\Lambda^*(x/t).$$

For

$$a > E \left[\frac{S_t^{(N)}}{N} \right] = \lambda t$$

the LDP result gives the assertion of the proposition (see [14, Lemma 2.6 and Theorem 2.8]). \square

APPENDIX B

PROOF OF RESULTS ON THE ASYMPTOTIC PROBABILITY OF DELAY VIOLATION

Lemma 2: Given $g \in \mathcal{G}$, $T \in \mathbb{T}$, $r > \lambda$, a batch service of rNT every T timeslots, and a g -smoothly scaling bit-arrival process characterized by the limiting g -scaled log moment generation function Λ , the decay rate of $P_{\text{delay}}(r, T)$ is given by the function I , i.e.

$$\lim_{N \rightarrow \infty} \frac{1}{g(N)} \log P_{\text{delay}}(r, T) = -I(r, T) \quad (36)$$

where

$$I(r, T) = \min_{\substack{t \in \mathbb{Z}^+ \\ tT + T - 1 - k > 0}} (tT + T - 1 - k) \times \Lambda^* \left(r + \frac{(D + 1 - 2T)r}{tT + T - 1 - k} \right) \quad (37)$$

and $k = D(\text{mod } T)$. In addition, $I(r, T)$ is lower-semicontinuous and increasing on r .

Proof: Let $g \in \mathcal{G}$, $T \in \mathbb{T} = \{1, 2, \dots, \lfloor \frac{D}{2} \rfloor\}$, $r > \lambda$, and $k = D(\text{mod } T)$. Without loss of generality, we assume that $I(r, T) < \infty$.

For any given SNR ρ and $N = \log \rho$, there are $A_t^{(N)}$ bits arriving at time t . The queue is being served exactly at times mT , for $m \in \mathbb{Z}$, with an instantaneous removal of the oldest $RT = rNT$ bits. The corresponding queue dynamics for the queue size $Q_t^{(N)}$, at time t , are as follows.

$$Q_t^{(N)} = \begin{cases} \left[Q_{t-1}^{(N)} + A_t^{(N)} - TR \right]^+, & \text{if } t = mT, m \in \mathbb{Z} \\ Q_{t-1}^{(N)} + A_t^{(N)}, & \text{otherwise} \end{cases} \quad (38)$$

where $Q_{-\infty}^{(N)} \equiv 0$. Since the arrival process is stationary and the system started empty at time $-\infty$, then $Q_i^{(N)}$ has the same steady-state distribution as that of $Q_{mT+i}^{(N)}$, $m \in \mathbb{Z}$, for each $i = 0, \dots, T-1$. The delay at time i also has the same steady-state distribution as the delay at time $mT+i$. Since $P_{\text{delay}}(r, T)$, as a function of r, T , is defined as the probability of the steady-state delay being greater than D , we have

$$\begin{aligned} P_{\text{delay}}(r, T) &: \\ &= P(\text{steady-state delay of a bit} > D) \\ &= \frac{1}{T} \sum_{i=0}^{T-1} P(\text{s-s delay of a bit arriving at time } i > D) \end{aligned} \quad (39)$$

where the equality holds since the arrivals are independent across time. From Lemma 4 in Appendix E, we have that the delay violation probability of *any* bit arriving at time i is

asymptotically equal to the delay violation probability of the last bit arriving at time i , (39) becomes

$$P_{\text{delay}}(r, T) \stackrel{a}{=} \frac{1}{T} \sum_{i=0}^{T-1} P(Q_i^{(N)}) \stackrel{a}{=} \sum_{i=0}^{T-1} P(Q_i^{(N)}) \quad (40)$$

where $Q_i^{(N)}$ denotes the event that the last bit arriving at timeslot i violates the delay bound D . This holds because T is a constant independent of ρ . Hence, (40) says that P_{delay} is asymptotically equal to the sum of $P(Q_i^{(N)})$.

Next, we relate the event $Q_i^{(N)}$ to a condition on the queue length $Q_i^{(N)}$, for $i = 0, \dots, T-1$. To do this, we need to describe the condition that the delay of the last bit arriving at timeslot i violates the delay bound D . Upon arrival, the last bit sees $Q_i^{(N)}$ bits (including itself) waiting in the queue. Since the batch service happens exactly in multiples of T , the bit must wait $T - i$ timeslots for the next service to start and another $\left\lceil \frac{Q_i^{(N)}}{RT} \right\rceil T$ timeslots for all $Q_i^{(N)}$ bits (including the last bit) to get served and be decoded. Hence, the last bit arriving at time i violates the delay bound D if, and only if

$$T - i + \left\lceil \frac{Q_i^{(N)}}{RT} \right\rceil T > D.$$

Let $\Omega^{(N)}$ contains all measurable random events. The condition above implies that the delay violation event for the last bit is given as

$$Q_i^{(N)} := \left\{ \omega \in \Omega^{(N)} : T - i + \left\lceil \frac{Q_i^{(N)}(\omega)}{RT} \right\rceil T > D \right\}. \quad (41)$$

Using (38) and (41), we show in Lemma 3 of Appendix D that

$$\begin{aligned} P_{\text{delay}}(r, T) &\stackrel{a}{=} P(Q_{T-1-k}^{(N)}) \\ &\stackrel{a}{=} P(Q_{T-1-k}^{(N)} > (D - T - k)R). \end{aligned} \quad (42)$$

Intuitively, this means that $P_{\text{delay}}(r, T)$ is asymptotically equal to $P(Q_{T-1-k}^{(N)})$, equivalently $P_{\text{delay}}(r, T)$ is asymptotically equal to the probability that the last bit arriving at time $T-1-k$ sees a queue length greater than $(D - T - k)R$ bits.

Finally, using (42), what remains is to establish that

$$\begin{aligned} &\lim_{N \rightarrow \infty} \frac{\log P(Q_{T-k-1}^{(N)} > (D - T - k)rN)}{g(N)} \\ &= -I(r, T) \\ &= - \min_{\substack{t \in \mathbb{Z}^+ \\ tT + T - 1 - k > 0}} (tT + T - k - 1) \\ &\quad \times \Lambda^* \left(r + \frac{(D + 1 - 2T)r}{tT + T - k - 1} \right). \end{aligned} \quad (43)$$

For notational simplicity, let $i := T - 1 - k$ and $q := (D - T - k)r$. Note that $q > ri \geq 0$ since $T \in \{1, 2, \dots, \lfloor \frac{D}{2} \rfloor\}$ and $k = D(\text{mod } T)$. Now, since

$$\frac{q + rTt}{Tt + i} = r + \frac{(D + 1 - 2T)r}{tT + T - k - 1}$$

it is sufficient to show that

$$\lim_{N \rightarrow \infty} \frac{\log P(Q_i^{(N)} > Nq)}{g(N)} = - \min_{\substack{t \in \mathbb{Z}^+ \\ tT+i > 0}} (Tt+i)\Lambda^* \left(\frac{q+rTt}{Tt+i} \right). \quad (44)$$

We separately show (matching) upper and lower bounds.

First, we show the lower bound. By using the queue dynamics in (38) recursively and the assumption of $Q_{-\infty}^{(N)} = 0$, the queue length $Q_i^{(N)}$ is related to the arrivals $A_j^{(N)}$, $j \leq i$, in the following manner:

$$Q_i^{(N)} = \sup_{t \in \mathbb{Z}^+} \left(\sum_{j=-tT+1}^i A_j^{(N)} - rtTN \right) \quad (45)$$

where we use the convention that $\sum_{j=1}^0 A_j^{(N)} = 0$. Using this relation and the fact that $q > 0$, we have

$$\begin{aligned} P(Q_i^{(N)} > Nq) &= P \left(\sup_{t \in \mathbb{Z}^+} \sum_{j=-tT+1}^i A_j^{(N)} - rtTN > Nq \right) \\ &= P \left(\sup_{\substack{t \in \mathbb{Z}^+ \\ tT+i > 0}} \sum_{j=-tT+1}^i A_j^{(N)} - rtTN > Nq \right). \end{aligned}$$

Now, for any fixed $t \in \mathbb{Z}^+$ so that $tT+i > 0$, we have

$$\begin{aligned} P(Q_i^{(N)} > Nq) &\geq P \left(\sum_{j=-tT+1}^i A_j^{(N)} - rtTN > Nq \right) \\ &= P \left(\sum_{j=1}^{tT+i} A_j^{(N)} > N(q+rTt) \right) \\ &= P \left(\frac{S_{Tt+i}^{(N)}}{N} > q+rTt \right). \end{aligned}$$

Taking the limit of both sides and using Proposition 1, we have

$$\liminf_{N \rightarrow \infty} \frac{\log P(Q_i^{(N)} > Nq)}{g(N)} \geq -(Tt+i)\Lambda^* \left(\frac{q+rTt}{Tt+i} \right). \quad (46)$$

Since t is arbitrary, maximizing the RHS over t gives the appropriate lower bound:

$$\begin{aligned} \liminf_{N \rightarrow \infty} \frac{\log P(Q_i^{(N)} > Nq)}{g(N)} &\geq - \inf_{\substack{t \in \mathbb{Z}^+ \\ tT+i > 0}} (Tt+i)\Lambda^* \left(\frac{q+rTt}{Tt+i} \right). \quad (47) \end{aligned}$$

For the upper bound, we use the following result from Lemma 5 in Appendix F:

$$\begin{aligned} \limsup_{N \rightarrow \infty} \frac{\log P(Q_i^{(N)} > Nq)}{g(N)} &\leq - \inf_{\substack{t \in \mathbb{Z}^+ \\ tT+i > 0}} (Tt+i)\Lambda^* \left(\frac{q+rTt}{Tt+i} \right) \end{aligned}$$

noting that the RHS is strictly greater than $-\infty$, by assumption. Hence, the lower and upper bounds coincide and (44) holds.

To complete the proof, we show the properties of $I(r, T)$ for $T \in \mathbb{T}$. First, I is increasing on $r \geq \lambda$ because $\Lambda^*(x)$ is increasing on $x \geq \lambda$ ([14, Lemma 2.7]). Second, $I(r, T)$ is lower semicontinuous on r because I is the minimum of a number of function Λ^* which are lower semicontinuous ([14, Lemma 2.7]). \square

Approximation 1: Relaxing the integer constraint in (13) gives the lower bound of I as

$$I(r, T) \geq \delta_r r(D+1-2T) =: I_{ir}(r, T) \quad (48)$$

where

$$\delta_r = \sup\{\theta > 0 : \Lambda(\theta) < \theta r\}. \quad (49)$$

Proof: By the definition of I , we have

$$\begin{aligned} I(r, T) &= \min_{\substack{t \in \mathbb{Z}^+ \\ tT+T-1-k > 0}} (tT+T-1-k) \\ &\quad \times \Lambda^* \left(r + r \frac{D-2T+1}{tT+T-1-k} \right) \\ &\geq \min_{\tau \in \mathbb{R}^+} \tau \Lambda^* \left(r + r \frac{D-2T+1}{\tau} \right) \\ &= \delta_r r(D-2T+1) \end{aligned}$$

where the last equality is a result of [14, Lemma 3.4] with δ_r defined as in (49). \square

APPENDIX C

PROOF OF THE MAIN RESULT

Proof of Theorem 1: Recall that

$$P_{\text{tot}}(r, T) := P_{\text{ch}}(r, T) + (1 - P_{\text{ch}}(r, T))P_{\text{delay}}(r, T) \quad (50)$$

where, from (1)

$$P_{\text{ch}}(r, T) \doteq \rho^{-d_{\text{ch}}(r, T)} \quad (51)$$

and, from Lemma 2

$$P_{\text{delay}}(r, T) \stackrel{g}{=} e^{-I(r, T)g(\log \rho)}. \quad (52)$$

Case 1: when $\lim_{N \rightarrow \infty} \frac{g(N)}{N} = \gamma \in (0, \infty)$. We have

$$P_{\text{delay}}(r, T) \doteq \rho^{-\gamma I(r, T)} \quad (53)$$

and

$$P_{\text{tot}}(r, T) \doteq \rho^{-\min\{\gamma I(r, T), d_{\text{ch}}(r, T)\}}. \quad (54)$$

The optimal negative SNR exponent of P_{tot} is

$$\begin{aligned} d^* &:= \sup_{\substack{r \in (\lambda, r_{\max}) \\ T \in \mathbb{T}}} \left\{ \lim_{\rho \rightarrow \infty} -\frac{\log P_{\text{tot}}(r, T)}{\log \rho} \right\} \\ &= \sup_{\substack{r \in (\lambda, r_{\max}) \\ T \in \mathbb{T}}} \{ \min \{ \gamma I(r, T), d_{\text{ch}}(r, T) \} \} \\ &= \max_{T \in \mathbb{T}} \left\{ \sup_{r \in (\lambda, r_{\max})} \{ \min \{ \gamma I(r, T), d_{\text{ch}}(r, T) \} \} \right\}. \end{aligned} \quad (55)$$

We first solve the optimization subproblem within the bracket for any given integer $T \in \mathbb{T}$. Because $I(r, T)$ is increasing on $r \geq \lambda$ while $d_{\text{ch}}(r, T)$ is strictly decreasing on $r \in [0, r_{\max}]$, the subproblem is solved by the optimal choice of multiplexing gain when the coding duration is fixed at T as

$$r^*(T) := \inf \{ r \in (\lambda, r_{\max}) : \gamma I(r, T) = d_{\text{ch}}(r, T) \}. \quad (56)$$

Hence, (55) is solved with the optimal coding duration T^* , given as

$$T^* = \arg \max_{T \in \mathbb{T}} \gamma I(r^*(T), T),$$

and the optimal multiplexing gain r^* , given as

$$r^* = r^*(T^*).$$

Note that, since $I(r, T) > 0$ when $r > \lambda$ and $d_{\text{ch}}(r, T) > 0$ when $r < r_{\max}$, it is guaranteed that $r^*(T) \in (\lambda, r_{\max})$.

Case 2: when

$$\lim_{N \rightarrow \infty} \frac{g(N)}{N} = 0 \quad \text{and} \quad \lim_{N \rightarrow \infty} \frac{g(N)}{\log N} = \infty.$$

In this case, for all $r \in (\lambda, r_{\max})$ and all $T \in \mathbb{T}$, we have $P_{\text{delay}}(r, T)$ asymptotically dominates $P_{\text{ch}}(r, T)$ and hence $P_{\text{tot}}(r, T)$ is asymptotically equal to $P_{\text{delay}}(r, T)$. Since, for any $T \in \mathbb{T}$, $I(r, T)$ is increasing on $r > \lambda$, we have

$$\begin{aligned} &\sup_{\substack{r \in (\lambda, r_{\max}) \\ T \in \mathbb{T}}} \lim_{\rho \rightarrow \infty} \frac{-\log P_{\text{tot}}(r, T)}{g(\log \rho)} \\ &\leq \max_{T \in \mathbb{T}} \left\{ \sup_{r \in (\lambda, r_{\max})} I(r, T) \right\} \\ &= \max_{T \in \mathbb{T}} I(r_{\max}, T). \end{aligned}$$

Case 3: when $\lim_{N \rightarrow \infty} \frac{g(N)}{N} = \infty$. This case is an opposite of Case 2. Here, $P_{\text{tot}}(r, T)$ is asymptotically equal to $P_{\text{ch}}(r, T)$ for

all $r \in (\lambda, r_{\max})$ and all $T \in \mathbb{T}$. Since $d_{\text{ch}}(r, T)$ is decreasing on r and increasing on T , we have

$$\begin{aligned} &\sup_{\substack{r \in (\lambda, r_{\max}) \\ T \in \mathbb{T}}} \lim_{\rho \rightarrow \infty} \frac{-\log P_{\text{tot}}(r, T)}{\log \rho} \\ &\leq \max_{T \in \mathbb{T}} \left\{ \sup_{r \in (\lambda, r_{\max})} d_{\text{ch}}(r, T) \right\} \\ &= \max_{T \in \mathbb{T}} d_{\text{ch}}(\lambda, T) \\ &= d_{\text{ch}} \left(\lambda, \left\lfloor \frac{D}{2} \right\rfloor \right). \quad \square \end{aligned}$$

APPENDIX D PROOF OF LEMMA 3

In this appendix, we prove the following lemma which is used in Appendix B.

Lemma 3: Consider $g \in \mathcal{G}$, $T \in \mathbb{T} = \{1, \dots, \lfloor \frac{D}{2} \rfloor\}$, $r > \lambda$, a family of g -smoothly scaling bit-arrival processes characterized by the limiting g -scaled log moment generation function Λ , and a periodic batch service of rNT bits at timeslots mT , $m \in \mathbb{Z}$. Let $Q_i^{(N)}$ be the queue length at time $i \in \{0, \dots, T-1\}$. Then, the event $\mathcal{Q}_{T-k-1}^{(N)}$, defined as

$$\mathcal{Q}_{T-k-1}^{(N)} = \left\{ \omega \in \Omega^{(N)} : k+1 + \left\lceil \frac{Q_{T-k-1}^{(N)}(\omega)}{RT} \right\rceil T > D \right\}$$

with $k = D(\text{mod } T)$, asymptotically dominates $P_{\text{delay}}(r, T)$. In other words

$$P_{\text{delay}}(r, T) \stackrel{g}{\sim} \Pr \left(\mathcal{Q}_{T-k-1}^{(N)} > (D - T - k)r \log \rho \right). \quad (57)$$

Proof: Let $k = D(\text{mod } T)$ and $i \in \{0, \dots, T-1\}$. Recall from (41) that

$$\mathcal{Q}_i^{(N)} = \left\{ \omega \in \Omega^{(N)} : T - i + \left\lceil \frac{Q_i^{(N)}(\omega)}{RT} \right\rceil T > D \right\}.$$

Now using the observation that, for any $x, y \in \mathbb{R}$

$$\lceil x \rceil > y \Leftrightarrow \lceil x \rceil > \lfloor y \rfloor \Leftrightarrow x > \lfloor y \rfloor$$

we have (58) shown at the bottom of the page.

$$\begin{aligned} \mathcal{Q}_i^{(N)} &= \left\{ \omega : \frac{Q_i^{(N)}(\omega)}{RT} > \left\lfloor \frac{D+i-T}{T} \right\rfloor \right\} \\ &= \begin{cases} \left\{ \omega : Q_i^{(N)}(\omega) > (D-T-k)R \right\}, & i \in [0, T-k-1] \\ \left\{ \omega : Q_i^{(N)}(\omega) > (D-k)R \right\}, & i \in [T-k, T-1]. \end{cases} \end{aligned} \quad (58)$$

On the other hand, (40) implies that

$$\begin{aligned}
P_{\text{delay}}(r, T) &\stackrel{\text{g}}{=} \sum_{i=0}^{T-1} P\left(\mathcal{Q}_i^{(N)}\right) \\
&= \sum_{i=0}^{T-k-1} P\left(\mathcal{Q}_i^{(N)}\right) + \sum_{i=T-k}^{T-1} P\left(\mathcal{Q}_i^{(N)}\right) \\
&\stackrel{\text{(a)}}{=} P\left(\mathcal{Q}_{T-k-1}^{(N)}\right) + P\left(\mathcal{Q}_{T-1}^{(N)}\right) \\
&\stackrel{\text{g}}{=} \max\left\{P\left(\mathcal{Q}_{T-k-1}^{(N)}\right), P\left(\mathcal{Q}_{T-1}^{(N)}\right)\right\} \\
&\stackrel{\text{(b)}}{=} \max\left\{P\left(\mathcal{Q}_{T-k-1}^{(N)} > (D-T-k)R\right), \right. \\
&\quad \left. P\left(\mathcal{Q}_{T-1}^{(N)} > (D-k)R\right)\right\} \\
&\stackrel{\text{(c)}}{=} P\left(\mathcal{Q}_{T-k-1}^{(N)} > (D-T-k)R\right) \quad (59)
\end{aligned}$$

where the equality in (b) is from (58). Next, we establish the (asymptotic) equalities (a) and (c). For (a), we first need to show that

$$\sum_{j=0}^{T-k-1} P\left(\mathcal{Q}_j^{(N)}\right) \stackrel{\text{g}}{=} P\left(\mathcal{Q}_{T-k-1}^{(N)}\right). \quad (60)$$

To establish this, we first observe that

$$\begin{aligned}
Q_j^{(N)}(\omega) &= Q_i^{(N)}(\omega) + \underbrace{A_{i+1}^{(N)}(\omega) + \dots + A_j^{(N)}(\omega)}_{\geq 0} \\
&\geq Q_i^{(N)}(\omega) \quad (61)
\end{aligned}$$

for all $\omega \in \Omega^{(N)}$ and $0 \leq i \leq j \leq T-1$. Hence, from (58), we have

$$P\left(\mathcal{Q}_{T-k-1}^{(N)}\right) \geq P\left(\mathcal{Q}_i^{(N)}\right), \quad i \in \{0, \dots, T-k-1\}$$

which implies

$$\sum_{i=0}^{T-k-1} P\left(\mathcal{Q}_i^{(N)}\right) \leq (T-k)P\left(\mathcal{Q}_{T-k-1}^{(N)}\right). \quad (62)$$

On the other hand, from the nonnegativity of probability, we have

$$\sum_{i=0}^{T-k-1} P\left(\mathcal{Q}_i^{(N)}\right) \geq P\left(\mathcal{Q}_{T-k-1}^{(N)}\right). \quad (63)$$

Combining (62) and (63), we have (60). Similarly, we can show that

$$\sum_{j=T-k}^{T-1} P\left(\mathcal{Q}_j^{(N)}\right) \stackrel{\text{g}}{=} P\left(\mathcal{Q}_{T-1}^{(N)}\right). \quad (64)$$

Combining (60) and (64), equality (a) in (59) is established.

To establish equality (c), it is sufficient to show that

$$\begin{aligned}
P\left(\mathcal{Q}_0^{(N)} > D'R\right) &\leq P\left(\mathcal{Q}_j^{(N)} > D'R\right) \\
&\leq P\left(\mathcal{Q}_0^{(N)} > (D'-T)R\right) \quad (65)
\end{aligned}$$

for any $D' > T$ and $j \in \{0, \dots, T-1\}$. This is because for $j_1 = T-1$ and $D'_1 = D-k$, we get

$$P\left(\mathcal{Q}_{T-1}^{(N)} > (D-k)R\right) \leq P\left(\mathcal{Q}_0^{(N)} > (D-T-k)R\right)$$

while for $j_2 = T-k-1$ and $D'_2 = D-T-k$, we get

$$P\left(\mathcal{Q}_0^{(N)} > (D-T-k)R\right) \leq P\left(\mathcal{Q}_{T-k-1}^{(N)} > (D-T-k)R\right)$$

asserting (c).

We prove (65) in two steps. The lower bound directly follows from (61), i.e.

$$Q_j^{(N)}(\omega) \geq Q_0^{(N)}(\omega), \quad \forall \omega \in \Omega^{(N)}.$$

For the upper bound, we notice that, for $D' > T$ and

$$\begin{aligned}
\omega \in \left\{\omega \in \Omega^{(N)} : Q_j^{(N)}(\omega) > D'R\right\} \\
\subseteq \left\{\omega \in \Omega^{(N)} : Q_j^{(N)}(\omega) > TR\right\}
\end{aligned}$$

$Q_j^{(N)}(\omega)$ is related to $Q_T^{(N)}(\omega)$ as

$$\begin{aligned}
Q_T^{(N)}(\omega) &= \left[Q_j^{(N)}(\omega) + A_{j+1}^{(N)}(\omega) + \dots + A_T^{(N)}(\omega) - TR\right]^+ \\
&= Q_j^{(N)}(\omega) + A_{j+1}^{(N)}(\omega) + \dots + A_T^{(N)}(\omega) - TR,
\end{aligned}$$

where $[\cdot]^+$ is removed. As a result, we have

$$\begin{aligned}
P\left(\mathcal{Q}_j^{(N)} > D'R\right) \\
&= P\left(\mathcal{Q}_T^{(N)} - \left\{A_{j+1}^{(N)} + \dots + A_T^{(N)}\right\} + TR > D'R\right) \\
&\leq P\left(\mathcal{Q}_T^{(N)} > (D'-T)R\right) \\
&= P\left(\mathcal{Q}_0^{(N)} > (D'-T)R\right)
\end{aligned}$$

where the last equality holds since $Q_T^{(N)}$ and $Q_0^{(N)}$ have the same stationary distribution. \square

APPENDIX E PROOF OF LEMMA 4

This appendix shows that the average probability of delay violation for bits that arrive at time i is asymptotically equal to the corresponding probability for the last bit arriving at that time. The proof is mainly based on the definition of the g -smoothly scaling process.

Lemma 4: Consider $g \in \mathcal{G}$ and a family of g -smoothly scaling bit-arrival processes $\left(\left(A_t^{(N)}, t \in \mathbb{Z}\right), N \in \mathbb{N}\right)$, characterized by the limiting g -scaled log moment generation function Λ . For any given N , let $W^{(N)}$ be a random variable having the same distribution as the steady-state distribution of the delay of a randomly chosen bit that arrives at time $i \in \{0, \dots, T-1\}$ while $Z^{(N)}$ is a random variable having a distribution that is

identical to the steady-state distribution of the delay for the last bit that arrives during time i . Then, for any $D > 0$

$$P(W^{(N)} > D) \stackrel{g}{=} P(Z^{(N)} > D). \quad (66)$$

Proof: We show (66) by showing the upper bound

$$P(W^{(N)} > D) \leq P(Z^{(N)} > D) \quad (67)$$

and the lower bound

$$P(W^{(N)} > D) \stackrel{g}{\geq} P(Z^{(N)} > D). \quad (68)$$

The upper bound is an immediate consequence of $W^{(N)}(\omega) \leq Z^{(N)}(\omega)$ for $\omega \in \Omega^{(N)}$. Below we prove the lower bound. We have

$$P(W^{(N)} > D) = \sum_{a \in \mathbb{N}} P(W^{(N)} > D | A_i^{(N)} = a) P(A_i^{(N)} = a). \quad (69)$$

Now, given that $A_i^{(N)} = a$ bits arrive at time i , we index the a bits as bit 1 to a , where bit 1 arrives first and bit a arrives last. Given $A_i^{(N)} = a$, we let $W_j^{(N)}$ to be the steady-state delay of the j th bit, $j \in \{1, \dots, a\}$. Since the bit can have any index, from 1 to a , with equal probability of $1/a$, we have

$$\begin{aligned} P(W^{(N)} > D | A_i^{(N)} = a) \\ = \frac{1}{a} \sum_{j=1}^a P(W_j^{(N)} > D | A_i^{(N)} = a). \end{aligned}$$

Ignoring all but the last term in the sum, we have

$$\begin{aligned} P(W^{(N)} > D | A_i^{(N)} = a) &\geq \frac{1}{a} P(W_a^{(N)} > D | A_i^{(N)} = a) \\ &= \frac{1}{a} P(Z^{(N)} > D | A_i^{(N)} = a) \end{aligned}$$

where the equality is a result of how $Z^{(N)}$ is defined. This means that

$$\begin{aligned} P(W^{(N)} > D) \\ &\geq \sum_{a \in \mathbb{N}} \frac{1}{a} P(Z^{(N)} > D | A_i^{(N)} = a) P(A_i^{(N)} = a) \\ &= \sum_{a \in \mathbb{N}} \frac{1}{a} P(Z^{(N)} > D \text{ and } A_i^{(N)} = a). \end{aligned}$$

Now, for a given $\beta > 0$, define

$$B^{(N)} := \{b \in \mathbb{N} : b < e^{\beta g(N)}\}.$$

We can further lower bound $P(W^{(N)} > D)$ as follows:

$$\begin{aligned} P(W^{(N)} > D) \\ &\geq \sum_{a \in B^{(N)}} \frac{1}{a} P(Z^{(N)} > D \text{ and } A_i^{(N)} = a) \\ &\geq e^{-\beta g(N)} \sum_{a \in B^{(N)}} P(Z^{(N)} > D \text{ and } A_i^{(N)} = a) \\ &= e^{-\beta g(N)} P(Z^{(N)} > D \text{ and } A_i^{(N)} \in B^{(N)}) \quad (70) \end{aligned}$$

where the second inequality holds because $1/a > e^{-\beta g(N)}$ for any $a \in B^{(N)}$.

Next, we show that $P(A_i^{(N)} \in B^{(N)}) \rightarrow 1$ as $N \rightarrow \infty$. We do this by using the definition of the g -smoothly scaling process: there exists $\theta > 0$ such that

$$\lim_{N \rightarrow \infty} \frac{\log E \left[e^{\theta A_i^{(N)} g(N)/N} \right]}{g(N)} = \Lambda(\theta) < \infty.$$

Hence, for any $\epsilon > 0$, there exists $N_0 = N_0(\epsilon)$ such that for all $N > N_0$, we have

$$g(N)(\Lambda(\theta) + \epsilon) > \log E \left[e^{\theta A_i^{(N)} g(N)/N} \right]. \quad (71)$$

The RHS can be lower bounded, for any $a_1 \in \mathbb{N}$:

$$\begin{aligned} \log E \left[e^{\theta A_i^{(N)} g(N)/N} \right] \\ &= \log \left(\sum_{a \in \mathbb{N}} P(A_i^{(N)} = a) e^{\theta a g(N)/N} \right) \\ &\geq \log \left(\sum_{a \geq a_1} P(A_i^{(N)} = a) e^{\theta a g(N)/N} \right) \\ &\geq \log \left(P(A_i^{(N)} \geq a_1) e^{\theta a_1 g(N)/N} \right) \\ &= \theta a_1 \frac{g(N)}{N} + \log P(A_i^{(N)} \geq a_1). \end{aligned}$$

This together with (71) gives

$$\log P(A_i^{(N)} \geq a_1) < g(N) \left[\Lambda(\theta) + \epsilon - \frac{\theta a_1}{N} \right]$$

for all $a_1 \in \mathbb{N}$. Now, we select $a_1 = e^{\beta g(N)}$ to get

$$\begin{aligned} \log \left(1 - P(A_i^{(N)} \in B^{(N)}) \right) &= \log P(A_i^{(N)} \geq e^{\beta g(N)}) \\ &< g(N) \left[\Lambda(\theta) + \epsilon - \frac{\theta e^{\beta g(N)}}{N} \right]. \end{aligned}$$

Since $\lim_{N \rightarrow \infty} \frac{g(N)}{\log N} = \infty$, we, then, have

$$P(A_i^{(N)} \in B^{(N)}) \rightarrow 1. \quad (72)$$

Finally, combining (72) and (70) implies that, for any $\beta > 0$,

$$\lim_{N \rightarrow \infty} \frac{\log P(W^{(N)} > D)}{g(N)} \geq \lim_{N \rightarrow \infty} \frac{\log P(Z^{(N)} > D)}{g(N)} - \beta.$$

Since β can be chosen arbitrarily small, we have the lower bound in (68), hence the assertion of the lemma. \square

APPENDIX F PROOF OF LEMMA 5

In this appendix, we prove the following lemma which is used in Appendix B.

Lemma 5: Consider $g \in \mathcal{G}$, $T \in \mathbb{T} = \{1, \dots, \lfloor \frac{D}{2} \rfloor\}$, $r > \lambda$, a family of g -smoothly scaling bit-arrival processes charac-

terized by the limiting g -scaled log moment generation function Λ , and a periodic batch service of rNT bits at timeslots mT , $m \in \mathbb{Z}$. Let $Q_i^{(N)}$ be the queue length at time $i \in \{0, \dots, T-1\}$. Then, for $q > ir$, we have

$$\limsup_{N \rightarrow \infty} \frac{\log P(Q_i^{(N)} > Nq)}{g(N)} \leq - \inf_{\substack{t \in \mathbb{Z}^+ \\ tT+i > 0}} (Tt+i)\Lambda^* \left(\frac{q+rTt}{Tt+i} \right) \quad (73)$$

assuming that the RHS is strictly greater than $-\infty$.

Proof: The proof uses the same technique as in [14, Lemmas 1.10 and 1.11]. Using (45), we have the following bound:

$$\begin{aligned} P(Q_i^{(N)} > Nq) &= P \left(\sup_{\substack{t \in \mathbb{Z}^+ \\ tT+i > 0}} \sum_{j=-tT+1}^i A_j^{(N)} - rtTN > Nq \right) \\ &= P \left(\sup_{\substack{t \in \mathbb{Z}^+ \\ tT+i > 0}} S_{tT+i}^{(N)} - rtTN > Nq \right) \\ &\leq \sum_{t > -\frac{i}{T}} P(S_{tT+i}^{(N)} > N(q+rTt)). \end{aligned}$$

Now, for any fixed $t_0 \in \mathbb{N}$, we have

$$\begin{aligned} P(Q_i^{(N)} > Nq) &\leq \sum_{-\frac{i}{T} < t \leq t_0} P(S_{tT+i}^{(N)} > N(q+rTt)) \\ &\quad + \sum_{t > t_0} P(S_{tT+i}^{(N)} > N(q+rTt)). \quad (74) \end{aligned}$$

Employing the principle of the largest term¹² gives

$$\begin{aligned} \limsup_{N \rightarrow \infty} \frac{\log P(Q_i^{(N)} > Nq)}{g(N)} &\leq \max \left(\max_{-\frac{i}{T} < t \leq t_0} \limsup_{N \rightarrow \infty} \frac{\log P(S_{tT+i}^{(N)} > N(q+rTt))}{g(N)}, \right. \\ &\quad \left. \limsup_{N \rightarrow \infty} \frac{1}{g(N)} \log \sum_{t > t_0} P(S_{tT+i}^{(N)} > N(q+rTt)) \right). \quad (75) \end{aligned}$$

¹²The principle of the largest term [14, Lemma 2.1]: Let a_n and b_n be sequences in \mathbb{R}^+ . If $n^{-1} \log a_n \rightarrow a$ and $n^{-1} \log b_n \rightarrow b$, then $n^{-1} \log(a_n + b_n) \rightarrow \max(a, b)$. This extends easily to finite sums.

For the first term (the $t \leq t_0$ term) in the maximum, we use Proposition 1 to get

$$\begin{aligned} &\max_{-\frac{i}{T} < t \leq t_0} \limsup_{N \rightarrow \infty} \frac{1}{g(N)} \log P \left(\frac{S_{tT+i}^{(N)}}{N} > q+rTt \right) \\ &\leq \max_{-\frac{i}{T} < t \leq t_0} -(Tt+i)\Lambda^* \left(\frac{q+rTt}{Tt+i} \right) \\ &\leq - \inf_{\substack{t \in \mathbb{Z}^+ \\ tT+i > 0}} (Tt+i)\Lambda^* \left(\frac{q+rTt}{Tt+i} \right) \quad (76) \end{aligned}$$

which is the RHS of (73) and finite by assumption.

Now, we show that we can select t_0 appropriately such that the second term (the $t > t_0$ term) in the RHS of (75) is also no greater than the RHS of (73). In other words, we show that there exists t_0 such that

$$\begin{aligned} \limsup_{N \rightarrow \infty} \frac{1}{g(N)} \log \sum_{t > t_0} P(S_{tT+i}^{(N)} > N(q+rTt)) \\ \leq - \inf_{\substack{t \in \mathbb{Z}^+ \\ tT+i > 0}} (Tt+i)\Lambda^* \left(\frac{q+rTt}{Tt+i} \right). \quad (77) \end{aligned}$$

This is shown by proving that there exist some $\theta > 0$ and $\epsilon > 0$ such that

$$\limsup_{N \rightarrow \infty} \frac{1}{g(N)} \log \sum_{t > t_0} P(S_{tT+i}^{(N)} > N(q+rTt)) \leq -\epsilon\theta((t_0+1)T+i) \quad (78)$$

for all $t_0 \in \mathbb{N}$. Now, selecting

$$t_0 = \left\lceil \frac{1}{\epsilon\theta T} \inf_{\substack{t \in \mathbb{Z}^+ \\ tT+i > 0}} (Tt+i)\Lambda^* \left(\frac{q+rTt}{Tt+i} \right) \right\rceil$$

provides (77).

To prove (78), we first use Chernoff bound as shown in (79) at the bottom of the page, where θ is an arbitrary positive scalar and the second equality is a consequence of i.i.d. assumption on $A_t^{(N)}$.

Next, we use the convexity of Λ and the fact that $\Lambda'(0) = \lambda < r$ (Remark 2) to establish that there exist some $\theta > 0$ and $\epsilon > 0$ for which

$$\Lambda(\theta) < \theta(r-2\epsilon). \quad (80)$$

$$\begin{aligned} \sum_{t > t_0} P(S_{tT+i}^{(N)} > N(q+rTt)) &= \sum_{t > t_0} P \left(e^{\frac{\theta g(N)}{N} S_{tT+i}^{(N)}} > e^{\frac{\theta g(N)}{N} N(q+rTt)} \right) \\ &\leq \sum_{t > t_0} e^{-\theta g(N)(q+rTt)} E \left[e^{\theta S_{tT+i}^{(N)} \frac{g(N)}{N}} \right] \\ &= \sum_{t > t_0} e^{-\theta g(N)(q+rTt)} \left(E \left[e^{\theta A_1^{(N)} \frac{g(N)}{N}} \right] \right)^{tT+i} \\ &= \sum_{t > t_0} \exp \left(-g(N)(tT+i) \left[\theta \left(\frac{q+rTt}{tT+i} \right) - \frac{\log E \left[e^{\theta A_1^{(N)} \frac{g(N)}{N}} \right]}{g(N)} \right] \right) \quad (79) \end{aligned}$$

On the other hand, from (5), we know that

$$\frac{\log E \left[e^{\frac{\theta g(N)}{N} A_1^{(N)}} \right]}{g(N)} \rightarrow \Lambda(\theta).$$

This means that there exists a $N_0 = N_0(\theta, \epsilon)$ such that, for all $N > N_0$

$$\frac{\log E \left[e^{\frac{\theta g(N)}{N} A_1^{(N)}} \right]}{g(N)} < \Lambda(\theta) + \theta\epsilon.$$

Combining this with (80), we have

$$\frac{\log E \left[e^{\frac{\theta g(N)}{N} A_1^{(N)}} \right]}{g(N)} < \theta(r - 2\epsilon) + \theta\epsilon = \theta(r - \epsilon), \quad (81)$$

for all $N > N_0$.

Hence, using (81), the term inside the square bracket in (79) can be bounded, uniformly over all $t > t_0$, as

$$\begin{aligned} & \theta \left(\frac{q + rtT}{tT + i} \right) - \frac{\log E \left[e^{\frac{\theta g(N)}{N} A_1^{(N)}} \right]}{g(N)} \\ &= \theta \left(r + \frac{q - ir}{tT + i} \right) - \frac{\log E \left[e^{\frac{\theta g(N)}{N} A_1^{(N)}} \right]}{g(N)} \\ &> \theta r - \frac{\log E \left[e^{\frac{\theta g(N)}{N} A_1^{(N)}} \right]}{g(N)} \\ &> \theta r - \theta(r - \epsilon) \\ &= \theta\epsilon \end{aligned} \quad (82)$$

where the first equality holds because $q > ir$, by assumption.

Inserting (82) into (79), we have (78)

$$\begin{aligned} & \limsup_{N \rightarrow \infty} \frac{1}{g(N)} \log \sum_{t > t_0} P \left(S_{tT+i}^{(N)} > N(q + rTt) \right) \\ & \leq \limsup_{N \rightarrow \infty} \frac{1}{g(N)} \log \sum_{t > t_0} \exp(-g(N)(tT + i)\theta\epsilon) \\ & = \limsup_{N \rightarrow \infty} \frac{1}{g(N)} \log \left(\frac{e^{-g(N)\theta\epsilon((t_0+1)T+i)}}{1 - e^{-g(N)\theta\epsilon T}} \right) \\ & = -\epsilon\theta((t_0 + 1)T + i) \end{aligned}$$

and, hence, the assertion of the lemma. \square

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