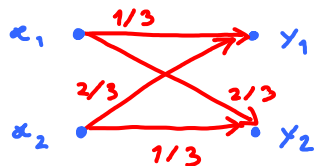


Q1

Wednesday, October 02, 2013 10:51 AM

(a) (i) The problem does not explicitly specify \mathcal{X} and \mathcal{Y} . However, from the size of the Q matrix, we know that $|\mathcal{X}| = |\mathcal{Y}| = 2$. So, we will set $\mathcal{X} = \{x_1, x_2\}$ and $\mathcal{Y} = \{y_1, y_2\}$.



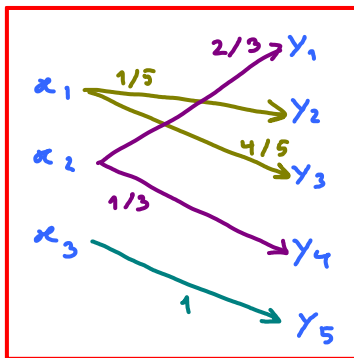
(ii) Note that this is a BSC with $p = \frac{2}{3}$. We know that the capacity of BSC is

$$C = 1 - H(p) = 1 - H\left(\frac{2}{3}\right) = 1 - 0.9183 \approx 0.0817$$

Alternatively, this channel is weakly symmetric. Therefore,

$$C = \log_2 |S_Y| - H(\underline{p}) = \log_2 2 - H\left[\left[\frac{1}{3} \quad \frac{2}{3}\right]\right] = 1 - 0.9183 \approx 0.0817$$

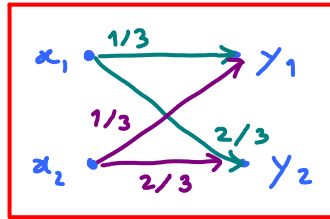
(b) (i) Let $\mathcal{X} = \{x_1, x_2, x_3\}$ and $\mathcal{Y} = \{y_1, y_2, y_3, y_4, y_5\}$.



(ii) Note that this is a noisy channel with non overlapping outputs. (Only one non-zero element in each column of Q .)

$$\text{So, } C = \log_2 |S_X| = \log_2 3 \approx 1.5850.$$

(c) (i) As in (a), we let $\mathcal{X} = \{x_1, x_2\}$ and $\mathcal{Y} = \{y_1, y_2\}$



(ii) The rows of Q are the same.

So, $Q(y|x)$ does not depend on x .

→ can "pull" $Q(y|x)$ out of the summation

$$\text{Therefore, } q(y) = \sum_x p(x) Q(y|x) = Q(y|x) \sum_x p(x) = Q(y|x),$$

which implies $X \perp\!\!\!\perp Y$.

Therefore, $I(X; Y) = 0$ for any input distribution.

Hence, $C = 0$.

Q2

Monday, October 27, 2014 4:51 PM

(a) There are $3 = |X|$ possible channel inputs and $4 = |Y|$ possible channel outputs. Therefore, the Q matrix must be 3×4 .

(b) We know that $C \leq \min \left\{ \log_2 |X|, \log_2 |Y| \right\} = \log_2 3 \approx 1.5850$
 $\log_2 3$ $\log_2 4$

However $1.6 > 1.5850$. Therefore $C = 1.6$ bpcu is impossible.

(c) Note that $\log_2 3 = \log_2 |X|$. Therefore,

let's try to construct a noiseless channel
 or a noisy channel with non-overlapping output

Examples of answers:



The correspond Q for the examples above are

$$Q = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix} \text{ or } Q = \begin{bmatrix} 1/2 & 1/2 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}, \text{ respectively.}$$

Q3a

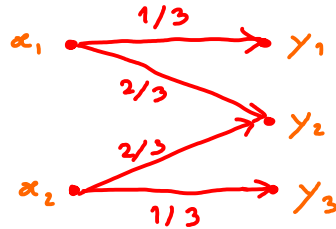
Monday, October 27, 2014 4:59 PM

The given DMC belong to a family of channels called Binary Erasure Channel (BEC). The general form of the channel diagram and the corresponding Q matrix are shown below:



Usually, $x_1 = y_1 = 0, x_2 = y_3 = 1, y_2 = e$
 ↑ the case where the bit is "erased"

Here, we have $\alpha = 2/3$. so, the channel diagram is



(b) For general α ,

$$Q = \begin{matrix} & \begin{matrix} 0 & e & 1 \end{matrix} \\ \begin{matrix} 0 \\ 1 \end{matrix} & \begin{bmatrix} 1-\alpha & \alpha & 0 \\ 0 & \alpha & 1-\alpha \end{bmatrix} \end{matrix} \left. \begin{array}{l} \rightarrow H(Y|X=0) = H(\alpha) \\ \rightarrow H(Y|X=1) = H(\alpha) \end{array} \right\} \Rightarrow H(Y|X) = H(\alpha)$$

Let $p = [p_0 \ 1-p_0]$. Then,

$$q_y = pQ = [p_0 \ 1-p_0] \begin{bmatrix} 1-\alpha & \alpha & 0 \\ 0 & \alpha & 1-\alpha \end{bmatrix} = \left[\underbrace{p_0(1-\alpha)}_{\text{write these as } \bar{a}} \quad \alpha \quad \underbrace{(1-p_0)(1-\alpha)}_{\text{write this as } \bar{p}_0} \right]$$

$$= [p_0 \bar{a} \quad \alpha \quad \bar{p}_0 \bar{a}]$$

Therefore, $H(Y) = -p_0 \bar{a} \log_2 p_0 \bar{a} - \alpha \log_2 \alpha - \bar{p}_0 \bar{a} \log_2 \bar{p}_0 \bar{a}$

and

$$I(X;Y) = H(Y) - H(Y|X) = H(Y) - H(\alpha).$$

Some useful facts: For any positive function $f(x)$,

$$\frac{d}{dx} f(x) \ln f(x) = f'(x) \ln f(x) + f(x) \frac{1}{f(x)} f'(x)$$

$$= f'(x) (1 + \ln f(x)) = f'(x) \ln(e f(x))$$

$$\left. \begin{array}{l} \times \frac{1}{\ln 2} \\ \downarrow \end{array} \right\} = f'(x) \log_2 e f(x)$$

$$\frac{d}{d\alpha} f(x) \log_2 f(x) = f'(x) \log_2 e f(x)$$

$$\frac{d}{dp_0} I(x;Y) = -\bar{\alpha} \log_2(\cancel{p_0} \bar{\alpha}) - (-\bar{\alpha}) \log_2(\cancel{e} \bar{p}_0 \bar{\alpha}) = 0$$

$$p_0 = \bar{p}_0$$

$$p_0 = 1 - p_0$$

$$p_0 = \frac{1}{2}$$

One may say that the uniform pmf is "obvious" from the "symmetry" in the transition probabilities from the inputs.

With $p_0 = \frac{1}{2}$, we have

$$C = I(x;Y) = -\frac{\bar{\alpha}}{2} \log_2 \frac{\bar{\alpha}}{2} - \cancel{\alpha} \log_2 \alpha - \frac{\bar{\alpha}}{2} \log_2 \frac{\bar{\alpha}}{2} - (-\cancel{\alpha} \log_2 \alpha - \bar{\alpha} \log_2 \bar{\alpha})$$

$$= -\bar{\alpha} \log_2 \frac{\bar{\alpha}}{2} + \bar{\alpha} \log_2 \bar{\alpha} = \bar{\alpha} \log_2 \frac{\bar{\alpha}}{\bar{\alpha}/2} = \bar{\alpha} = 1 - \alpha$$

When $\alpha = \frac{2}{3}$, $C = 1 - \frac{2}{3} = \frac{1}{3}$ bpcu.

The capacity-achieving \bar{p} is $\bar{p} = \left[\frac{1}{2} \quad \frac{1}{2} \right]$.

$$\text{So, } I(X; Y) = H(Y) - H(Y|X) = -\frac{1+p_1}{2} \log_2 \frac{1+p_1}{2} - \frac{1-p_1}{2} \log_2 \frac{1-p_1}{2} - (1-p_1)$$

$$\frac{d}{dp_1} I(X; Y) = -\frac{1}{2} \log_2 e \frac{1+p_1}{2} - \left(-\frac{1}{2}\right) \log_2 e \frac{1-p_1}{2} + 1 = 0$$

$$\frac{1}{2} \log_2 \frac{1-p_1}{1+p_1} + 1 = 0$$

$$\frac{1+p_1}{1-p_1} = 2^2 = 4$$

$$1+p_1 = 4 - 4p_1$$

$$5p_1 = 3$$

$$p_1 = \frac{3}{5}$$

With $p_1 = \frac{3}{5}$,

$$\mathcal{P} = \left[\frac{1+\frac{3}{5}}{2} \quad \frac{1-\frac{3}{5}}{2} \right] = \left[\frac{4}{5} \quad \frac{1}{5} \right]$$

$$\text{So, } C = H\left(\left[\frac{4}{5} \quad \frac{1}{5}\right]\right) - \left(1 - \frac{3}{5}\right) \approx 0.7219 - \frac{2}{5} = 0.3219 \text{ bpcu}$$

The pmf p of X that achieves this capacity value is

$$p = \left[\frac{3}{5} \quad p_2 \quad p_3 \right] \quad \text{where } p_2 + p_3 = \frac{2}{5}$$

↑

$$\text{Ex. } p_2 = p_3 = \frac{1}{5}$$

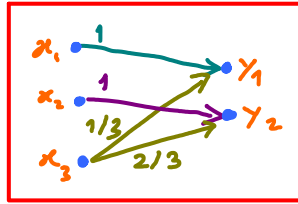
Q4

Monday, October 27, 2014 4:30 PM

(a)

$$Q = \begin{matrix} & \begin{matrix} y_1 & y_2 \end{matrix} \\ \begin{matrix} x_1 \\ x_2 \\ x_3 \end{matrix} & \begin{bmatrix} 1 & 0 \\ 0 & 1 \\ 1/3 & 2/3 \end{bmatrix} \end{matrix}$$

\Rightarrow



(b)

Note that $|Y| = 2$. So, $I(X; Y) \leq H(Y) \leq \log_2 |Y| = \log_2 2 = 1$.

Note also that without the last row, this Q corresponds to a noiseless channel with capacity $\log_2 |X| = \log_2 2 = 1$ which is achieved by uniform pmf on the input. ($p = [\frac{1}{2} \frac{1}{2}]$.)

Now, when the last row of Q is included, we may choose not to use it by using $p = [\frac{1}{2} \frac{1}{2} 0]$. This gives

$$p = [\frac{1}{2} \frac{1}{2} 0] \begin{bmatrix} 1 & 0 \\ 0 & 1 \\ 1/3 & 2/3 \end{bmatrix} = [\frac{1}{2} \frac{1}{2}]$$

so, $H(Y) = H([\frac{1}{2} \frac{1}{2}]) = 1$ and

$$H(Y|X) = \frac{1}{2} \times H([1 \ 0]) + \frac{1}{2} \times H([0 \ 1]) + 0 \times H([\frac{1}{3} \ \frac{2}{3}]) = 0.$$

Therefore, $I(X, Y) = H(Y) - H(Y|X) = 1 - 0 = 1$ which is the same as the bound above.

Because $I(X; Y)$ can not exceed the bound, we know that this $p = [\frac{1}{2} \frac{1}{2} 0]$ achieved the maximum $I(X; Y)$.

Therefore, $C = 1$ bpcu and corresponding $p = [\frac{1}{2} \frac{1}{2} 0]$

Q5

Friday, October 31, 2014 9:53 AM

close all; clear all;

<pre>%% 1.a Q = [1/3 2/3; 2/3 1/3]; [ps C] = capacity_blahut(Q)</pre>		<pre>>> Capacity_HW_blahut ps = [0.5000 0.5000] = p C = 0.0817 bpcu</pre>
<pre>%% 1.b Q = [0 1/5 4/5 0 0; 2/3 0 0 1/3 0; 0 0 0 0 1]; [ps C] = capacity_blahut(Q)</pre>		<pre>ps = [0.3333 0.3333 0.3333] = p C = 1.5850 bpcu</pre>
<pre>%% 1.c Q = [1/3 2/3; 1/3 2/3]; [ps C] = capacity_blahut(Q)</pre>		<pre>ps = [0.5000 0.5000] = p C = 0 bpcu</pre>
<pre>%% 2.c Q = [1 0 0 0; 0 1 0 0; 0 0 1 0]; [ps C] = capacity_blahut(Q)</pre>		<pre>ps = [0.3333 0.3333 0.3333] = p C = 1.5850 bpcu</pre>
<pre>%% 3.a Q = [1/3 2/3 0; 0 2/3 1/3]; [ps C] = capacity_blahut(Q)</pre>		<pre>ps = [0.5000 0.5000] = p C = 0.3333</pre> <p style="text-align: right; color: magenta;"><i>one of the answers</i></p>
<pre>%% 3.b Q = [1 0; 1/2 1/2; 1/2 1/2]; [ps C] = capacity_blahut(Q)</pre>		<pre>ps = [0.6000 0.2000 0.2000] = p C = 0.3219 bpcu</pre>
<pre>%% 4 Q = [1 0; 0 1; 1/3 2/3]; [ps C] = capacity_blahut(Q)</pre>		<pre>ps = [0.5000 0.5000 0.0000] = p C = 1.0000 bpcu</pre>